

Optimierende Compiler

Überblick über die GNU Compiler Collection – gcc

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FG Eingebettete Systeme und ihre Anwendungen
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Sommersemester 2011

- Nicht mehr nur GNU C Compiler
 - C, C++, Java, Ada, Objective-C, Objective-C++, Fortran
 - Go, Modula-2, Modula-3, Pascal, PL/I, Mercury, VHDL
 - UPC

Ab hier Material von Prof. Uday Khedker aus

Essential Abstractions in GCC '11

A Workshop on GCC Internals by GCC Resource Center
Department of Computer Science & Engineering Indian
Institute of Technology, Bombay

<http://www.cse.iitb.ac.in/grc/gcc-workshop-11/>

OC

A. Koch

Compile-Fluß

Aufbau

GIMPLE IR

RTL IR

Optimierung

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Part 4

*$GCC \equiv$ The **G**reat **C**ompiler **C**hallenge*

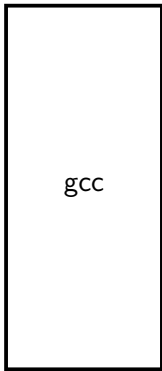
What is GCC?

- For the GCC developer community: [The GNU Compiler Collection](#)
- For other compiler writers: [The Great Compiler Challenge](#) 😊



The GNU Tool Chain for C

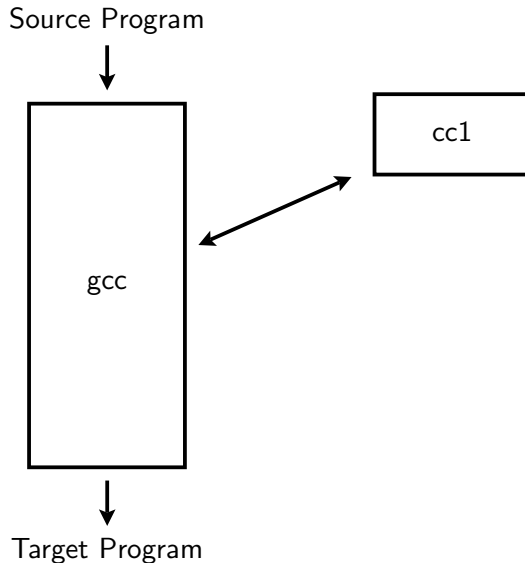
Source Program



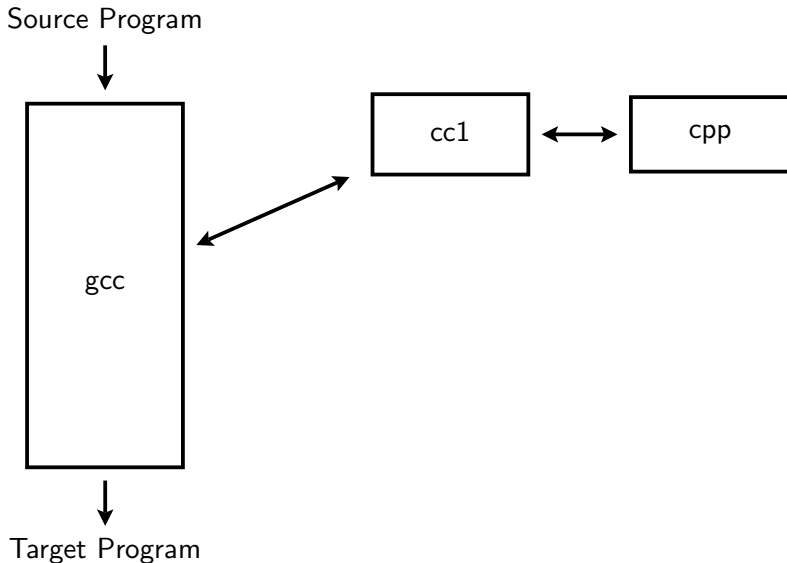
Target Program



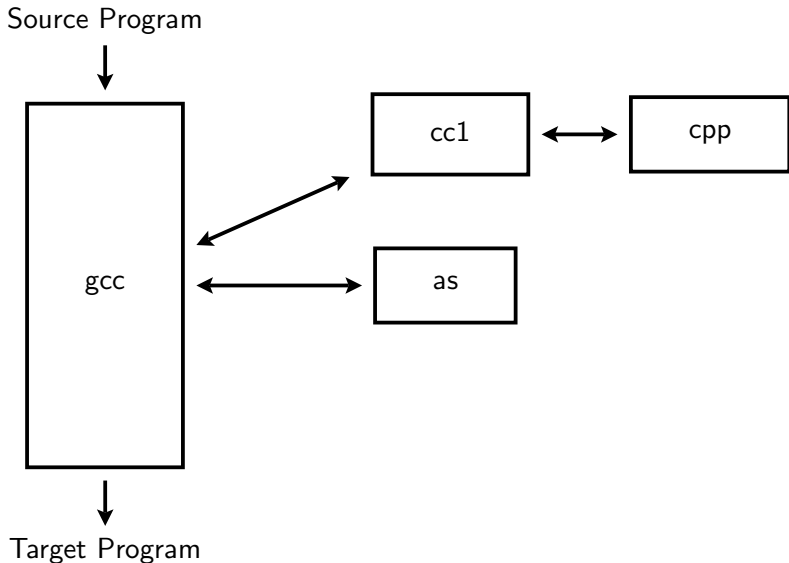
The GNU Tool Chain for C



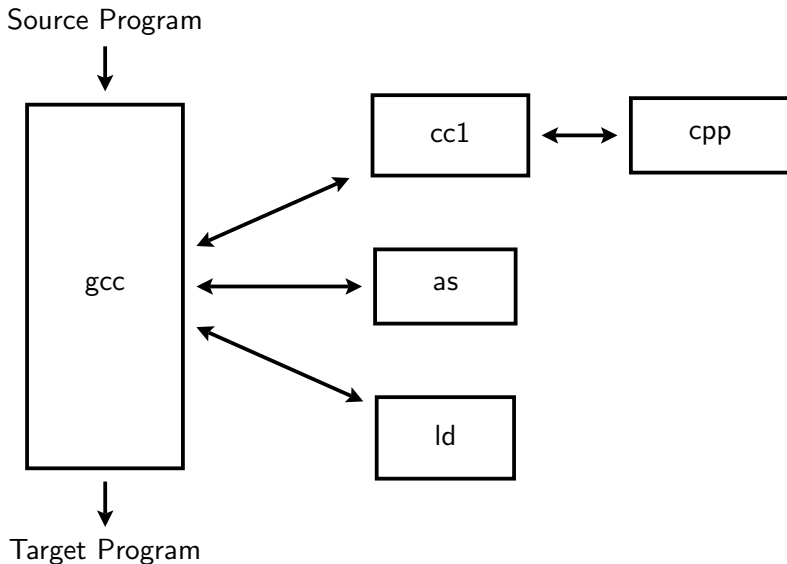
The GNU Tool Chain for C



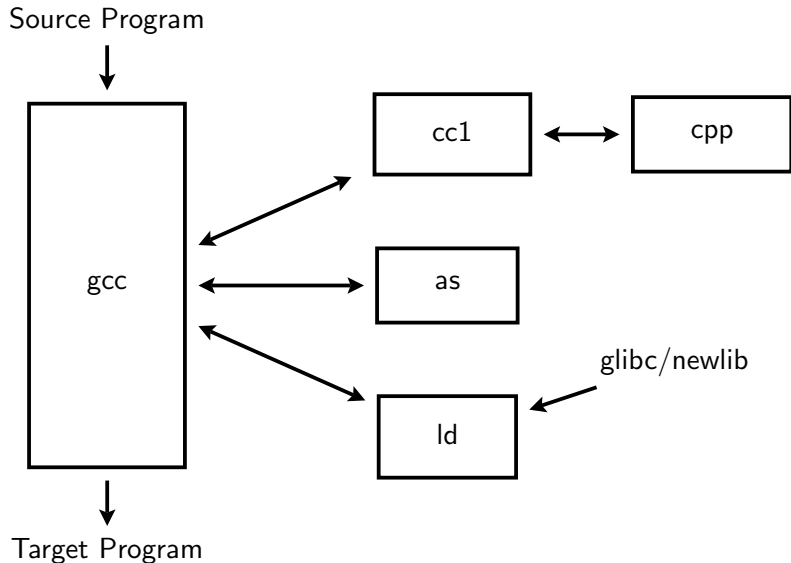
The GNU Tool Chain for C



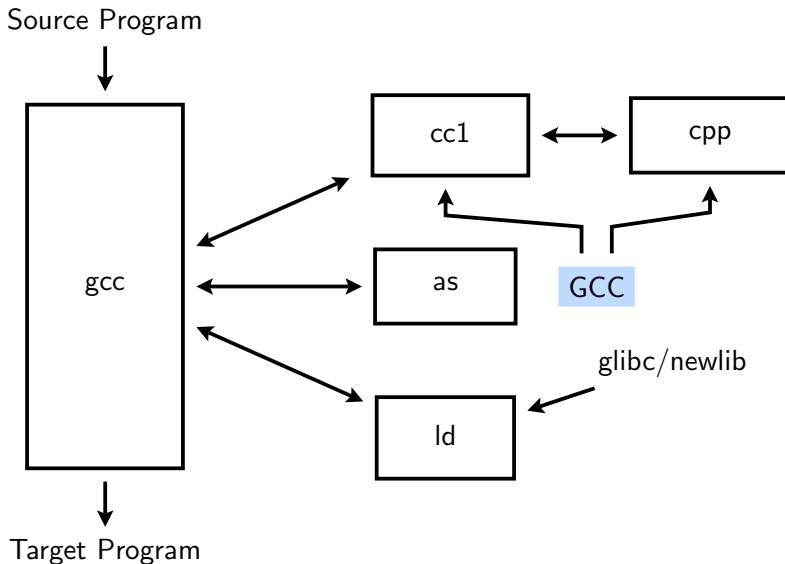
The GNU Tool Chain for C



The GNU Tool Chain for C



The GNU Tool Chain for C



Aufbau des GCC

Why is Understanding GCC Difficult?

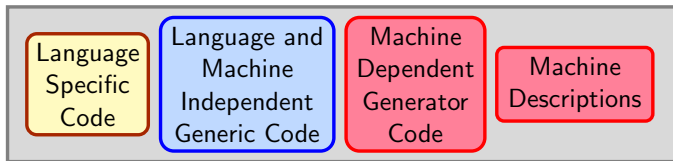
Deeper technical reasons

- GCC is not a compiler but a *compiler generation framework*
Two distinct gaps that need to be bridged
 - ▶ Input-output of the generation framework
The target specification and the generated compiler
 - ▶ Input-output of the generated compiler
A source program and the generated assembly program
- GCC generated compiler uses a derivative of the Davidson-Fraser model of compilation
 - ▶ Early instruction selection
 - ▶ Machine dependent intermediate representation
 - ▶ Simplistic instruction selection and retargetability mechanism



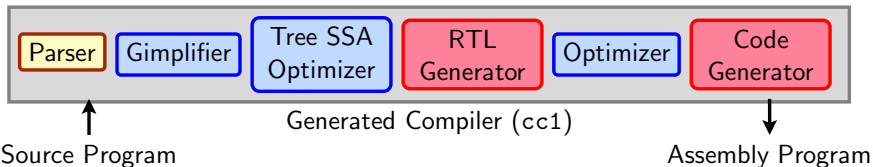
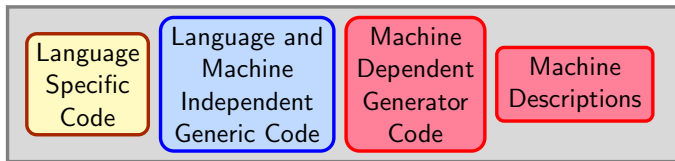
The Architecture of GCC

Compiler Generation Framework

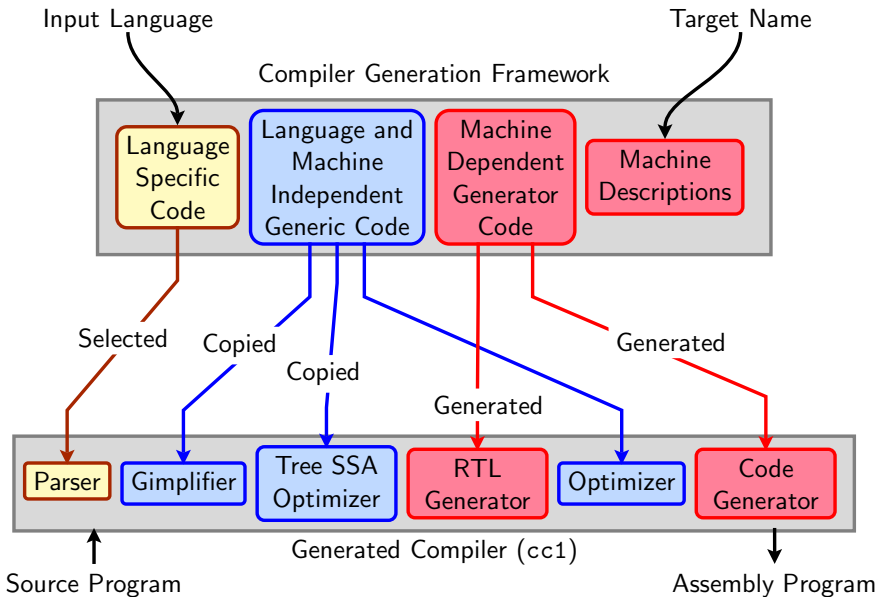


The Architecture of GCC

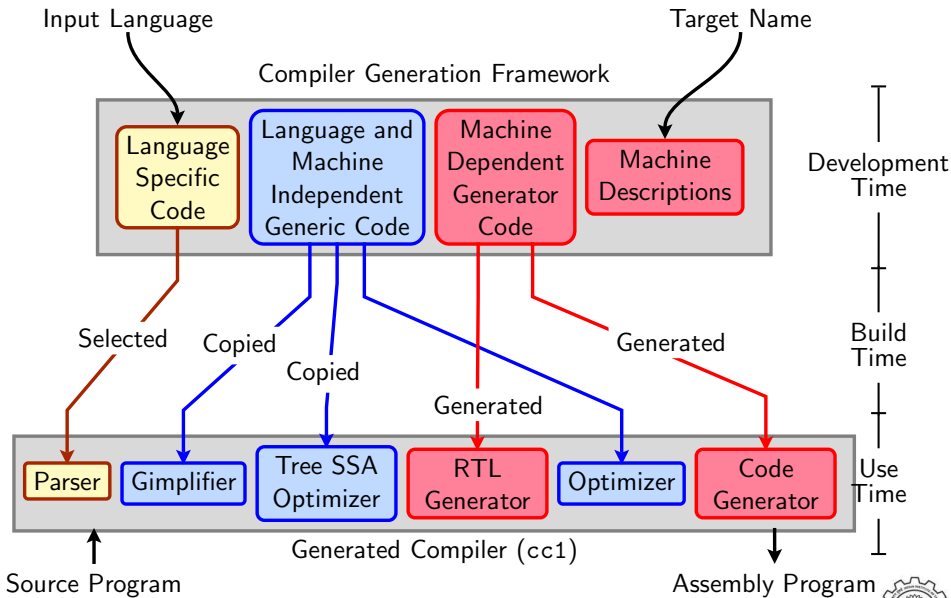
Compiler Generation Framework



The Architecture of GCC



The Architecture of GCC



OC

A. Koch

Compile-Fluß

Aufbau

GIMPLE IR

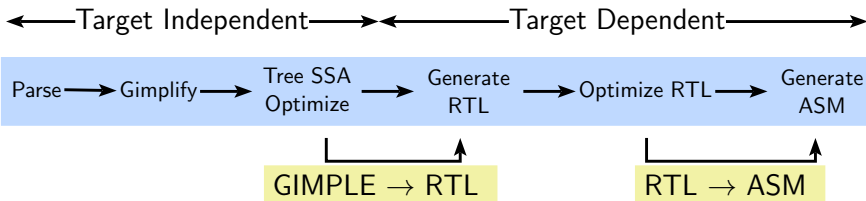
RTL IR

Optimierung

Zwischendarstellung GIMPLE

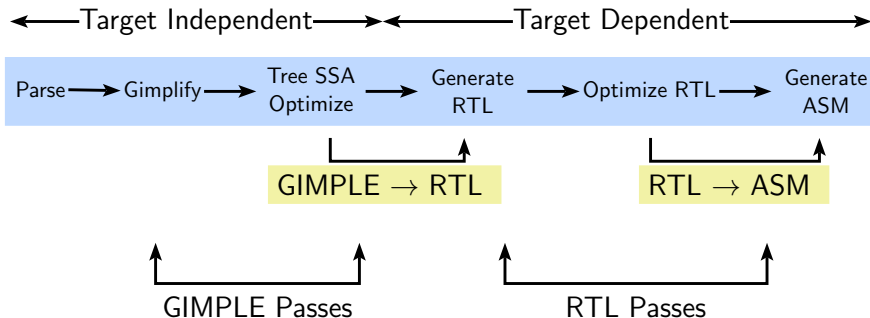
Basic Transformations in GCC

Transformation from a language to a *different* language



Basic Transformations in GCC

Transformation from a language to a *different* language



Transformation Passes in GCC 4.6.0

- A total of 207 unique pass names initialized in `$(SOURCE)/gcc/passes.c`
Total number of passes is 241.
 - ▶ Some passes are called multiple times in different contexts
Conditional constant propagation and dead code elimination are called thrice
 - ▶ Some passes are enabled for specific architectures
 - ▶ Some passes have many variations (eg. special cases for loops)
Common subexpression elimination, dead code elimination
- The pass sequence can be divided broadly in two parts
 - ▶ Passes on GIMPLE
 - ▶ Passes on RTL
- Some passes are organizational passes to group related passes



Passes On GIMPLE in GCC 4.6.0

Pass Group	Examples	Number of passes
Lowering	GIMPLE IR, CFG Construction	10
Simple Interprocedural Passes (Non-LTO)	Conditional Constant Propagation, Inlining, SSA Construction	38
Regular Interprocedural Passes (LTO)	Constant Propagation, Inlining, Pointer Analysis	10
LTO generation passes		02
Other Intraprocedural Optimizations	Constant Propagation, Dead Code Elimination, PRE Value Range Propagation, Rename SSA	65
Loop Optimizations	Vectorization, Parallelization, Copy Propagation, Dead Code Elimination	28
Generating RTL		01
<i>Total number of passes on GIMPLE</i>		154



Passes On RTL in GCC 4.6.0

Pass Group	Examples	Number of passes
Intraprocedural Optimizations	CSE, Jump Optimization, Dead Code Elimination, Jump Optimization	27
Loop Optimizations	Loop Invariant Movement, Peeling, Unswitching	07
Machine Dependent Optimizations	Register Allocation, Instruction Scheduling, Peephole Optimizations	50
Assembly Emission and Finishing		03
<i>Total number of passes on RTL</i>		87



Finding Out List of Optimizations

Along with the associated flags

- A complete list of optimizations with a brief description

```
gcc -c --help=optimizers
```

- Optimizations enabled at level 2 (other levels are 0, 1, 3, and s)

```
gcc -c -O2 --help=optimizers -Q
```



Producing the Output of GCC Passes

- Use the option `-fdump-<ir>-<passname>`

`<ir>` could be

- ▶ `tree`: Intraprocedural passes on GIMPLE
- ▶ `ipa`: Interprocedural passes on GIMPLE
- ▶ `rtl`: Intraprocedural passes on RTL

- Use `all` in place of `<pass>` to see all dumps

Example: `gcc -fdump-tree-all -fdump-rtl-all test.c`

- Dumping more details:

Suffix `raw` for tree passes and `details` or `slim` for RTL passes

Individual passes may have more verbosity options (e.g.

`-fsched-verbose=5`)

- Use `-S` to stop the compilation with assembly generation
- Use `--verbose-asm` to see more detailed assembly dump



Total Number of Dumps

Optimization Level	Number of Dumps	Goals
Default	47	Fast compilation
O1	134	
O2	158	
O3	168	
Os	156	Optimize for space



Selected Dumps for Our Example Program

GIMPLE dumps (t)

001t.tu

003t.original

004t.gimple

006t.vcg

009t.omplower

010t.lower

012t.eh

013t.cfg

017t.ssa

018t.veclower

019t.inline_param1

020t.einline

037t.release_ssa

038t.inline_param2

044i.whole-program

048i.inline

138t.cplxlower0

143t.optimized

224t.statistics

ipa dumps (i)

000i.cgraph

014i.visibility

015i.early_local_cleanups

044i.whole-program

048i.inline

rtl dumps (r)

144r.expand

145r.sibling

147r.initvals

148r.unshare

149r.vregs

150r.into_cfglayout

151r.jump

163r.reginfo

183r.outof_cfglayout

184r.split1

186r.dfinit

187r.mode_sw

188r.asmcons

191r.ira

194r.split2

198r.pro_and_epilogue

211r.stack

212r.alignments

215r.mach

216r.barriers

220r.shorten

221r.nothrow

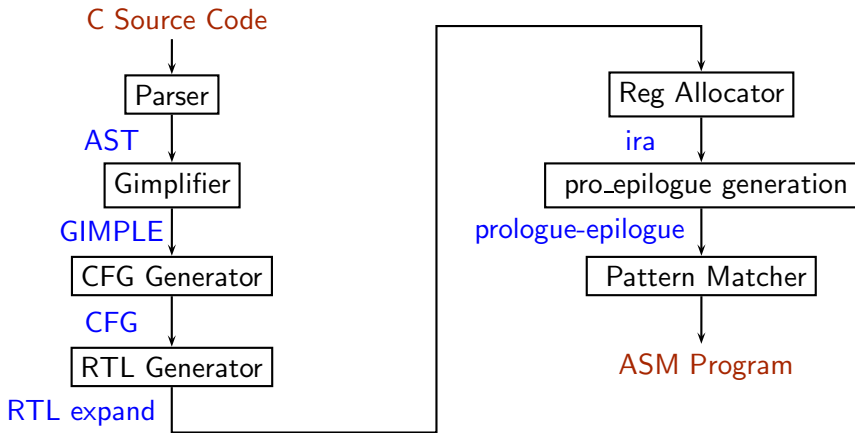
222r.final

223r.dfinish

assembly



Passes for First Level Graybox Probing of GCC



Lowering of abstraction!



Part 2

Examining GIMPLE Dumps

Gimplifier

- About GIMPLE
 - ▶ Three-address representation derived from GENERIC
Computation represented as a sequence of basic operations
Temporaries introduced to hold intermediate values
 - ▶ Control construct are explicated into conditional jumps
- Examining GIMPLE Dumps
 - ▶ Examining translation of data accesses
 - ▶ Examining translation of control flow
 - ▶ Examining translation of function calls



GIMPLE: Composite Expressions Involving Local and Global Variables

test.c

```
int a;

int main()
{
    int x = 10;
    int y = 5;

    x = a + x * y;
    y = y - a * x;
}
```

test.c.004t.gimple

```
x = 10;
y = 5;
D.1954 = x * y;
a.0 = a;
x = D.1954 + a.0;
a.1 = a;
D.1957 = a.1 * x;
y = y - D.1957;
```

Global variables are treated as “memory locations” and local variables are treated as “registers”



GIMPLE: Composite Expressions Involving Local and Global Variables

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  int x = 10;
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GIMPLE: Composite Expressions Involving Local and Global Variables

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test.c
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int a;
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int main()
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{
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  int x = 10;
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```

```
  x = a + x * y;
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```
  y = y - a * x;
```

```
}
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```
test.c.004t.gimple
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x = 10;
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D.1954 = x * y;
```

```
a.0 = a;
```

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x = D.1954 + a.0;
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a.1 = a;
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D.1957 = a.1 * x;
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GIMPLE: Composite Expressions Involving Local and Global Variables

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  y = y - a * x;
```

```
}
```

```
test.c.004t.gimple
```

```
x = 10;
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D.1954 = x * y;
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a.0 = a;
```

```
x = D.1954 + a.0;
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```
a.1 = a;
```

```
D.1957 = a.1 * x;
```

```
y = y - D.1957;
```

Global variables are treated as “memory locations” and local variables are treated as “registers”



GIMPLE: 1-D Array Accesses

test.c

```
int main()
{
    int a[3], x;
    a[1] = a[2] = 10;
    x = a[1] + a[2];
    a[0] = a[1] + a[1]*x;
}
```

test.c.004t.gimple

```
a[2] = 10;
D.1952 = a[2];
a[1] = D.1952;
D.1953 = a[1];
D.1954 = a[2];
x = D.1953 + D.1954;
D.1955 = x + 1;
D.1956 = a[1];
D.1957 = D.1955 * D.1956;
a[0] = D.1957;
```



GIMPLE: 1-D Array Accesses

test.c

```
int main()
{
    int a[3], x;
    a[1] = a[2] = 10;
    x = a[1] + a[2];
    a[0] = a[1] + a[1]*x;
}
```

test.c.004t.gimple

```
a[2] = 10;
D.1952 = a[2];
a[1] = D.1952;
D.1953 = a[1];
D.1954 = a[2];
x = D.1953 + D.1954;
D.1955 = x + 1;
D.1956 = a[1];
D.1957 = D.1955 * D.1956;
a[0] = D.1957;
```



GIMPLE: 1-D Array Accesses

test.c

```
int main()
{
    int a[3], x;
    a[1] = a[2] = 10;
    x = a[1] + a[2];
    a[0] = a[1] + a[1]*x;
}
```

test.c.004t.gimple

```
a[2] = 10;
D.1952 = a[2];
a[1] = D.1952;
D.1953 = a[1];
D.1954 = a[2];
x = D.1953 + D.1954;
D.1955 = x + 1;
D.1956 = a[1];
D.1957 = D.1955 * D.1956;
a[0] = D.1957;
```



GIMPLE: 1-D Array Accesses

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int main()
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    int a[3], x;
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    x = a[1] + a[2];
    a[0] = a[1] + a[1]*x;
}
```

test.c.004t.gimple

```
a[2] = 10;
D.1952 = a[2];
a[1] = D.1952;
D.1953 = a[1];
D.1954 = a[2];
x = D.1953 + D.1954;
D.1955 = x + 1;
D.1956 = a[1];
D.1957 = D.1955 * D.1956;
a[0] = D.1957;
```



GIMPLE: 1-D Array Accesses

test.c

```
int main()
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    int a[3], x;
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    a[0] = a[1] + a[1]*x;
}
```

test.c.004t.gimple

```
a[2] = 10;
D.1952 = a[2];
a[1] = D.1952;
D.1953 = a[1];
D.1954 = a[2];
x = D.1953 + D.1954;
D.1955 = x + 1;
D.1956 = a[1];
D.1957 = D.1955 * D.1956;
a[0] = D.1957;
```



GIMPLE: 2-D Array Accesses

test.c

```
int main()
{
    int a[3][3], x, y;
    a[0][0] = 7;
    a[1][1] = 8;
    a[2][2] = 9;
    x = a[0][0] / a[1][1];
    y = a[1][1] % a[2][2];
}
```

test.c.004t.gimple

```
a[0][0] = 7;
a[1][1] = 8;
a[2][2] = 9;
D.1953 = a[0][0];
D.1954 = a[1][1];
x = D.1953 / D.1954;
D.1955 = a[1][1];
D.1956 = a[2][2];
y = D.1955 % D.1956;
```



GIMPLE: 2-D Array Accesses

test.c

```
int main()
{
    int a[3][3], x, y;
    a[0][0] = 7;
    a[1][1] = 8;
    a[2][2] = 9;
    x = a[0][0] / a[1][1];
    y = a[1][1] % a[2][2];
}
```

test.c.004t.gimple

```
a[0][0] = 7;
a[1][1] = 8;
a[2][2] = 9;
D.1953 = a[0][0];
D.1954 = a[1][1];
x = D.1953 / D.1954;
D.1955 = a[1][1];
D.1956 = a[2][2];
y = D.1955 % D.1956;
```

- No notion of “addressable memory” in GIMPLE.
- Array reference is a single operation in GIMPLE and is linearized in RTL during expansion



GIMPLE: Use of Pointers

test.c

```
int main()
{
    int **a,*b,c;
    b = &c;
    a = &b;
    **a = 10; /* c = 10 */
}
```

test.c.004t.gimple

```
main ()
{
    int * D.1953;
    int * * a;
    int * b;
    int c;

    b = &c;
    a = &b;
    D.1953 = *a;
    *D.1953 = 10;
}
```



GIMPLE: Use of Pointers

test.c

```
int main()
{
    int **a,*b,c;
    b = &c;
    a = &b;
    **a = 10; /* c = 10 */
}
```

test.c.004t.gimple

```
main ()
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    int * D.1953;
    int * * a;
    int * b;
    int c;

    b = &c;
    a = &b;
    D.1953 = *a;
    *D.1953 = 10;
}
```



GIMPLE: Use of Structures

test.c

```
typedef struct address
{ char *name;
} ad;

typedef struct student
{ int roll;
  ad *ct;
} st;

int main()
{ st *s;
  s = malloc(sizeof(st));
  s->roll = 1;
  s->ct=malloc(sizeof(ad));
  s->ct->name = "Mumbai";
}
```

test.c.004t.gimple

```
main ()
{
  void * D.1957;
  struct ad * D.1958;
  struct st * s;
  extern void * malloc (unsigned int);

  s = malloc (8);
  s->roll = 1;
  D.1957 = malloc (4);
  s->ct = D.1957;
  D.1958 = s->ct;
  D.1958->name = "Mumbai";
}
```



GIMPLE: Use of Structures

test.c

```
typedef struct address
{ char *name;
} ad;

typedef struct student
{ int roll;
  ad *ct;
} st;

int main()
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  s = malloc(sizeof(st));
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test.c.004t.gimple

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  s = malloc (8);
  s->roll = 1;
  D.1957 = malloc (4);
  s->ct = D.1957;
  D.1958 = s->ct;
  D.1958->name = "Mumbai";
}
```



GIMPLE: Use of Structures

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typedef struct address
{ char *name;
} ad;

typedef struct student
{ int roll;
  ad *ct;
} st;

int main()
{ st *s;
  s = malloc(sizeof(st));
  s->roll = 1;
  s->ct=malloc(sizeof(ad));
  s->ct->name = "Mumbai";
}
```

test.c.004t.gimple

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main ()
{
  void * D.1957;
  struct ad * D.1958;
  struct st * s;
  extern void * malloc (unsigned int);

  s = malloc (8);
  s->roll = 1;
  D.1957 = malloc (4);
  s->ct = D.1957;
  D.1958 = s->ct;
  D.1958->name = "Mumbai";
}
```



GIMPLE: Use of Structures

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int main()
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  s = malloc(sizeof(st));
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  s->ct->name = "Mumbai";
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  void * D.1957;
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  struct st * s;
  extern void * malloc (unsigned int);

  s = malloc (8);
  s->roll = 1;
  D.1957 = malloc (4);
  s->ct = D.1957;
  D.1958 = s->ct;
  D.1958->name = "Mumbai";
}
```



GIMPLE: Pointer to Array

test.c

```
int main()
{
    int *p_a, a[3];

    p_a = &a[0];

    *p_a = 10;
    *(p_a+1) = 20;
    *(p_a+2) = 30;
}
```

test.c.004t.gimple

```
main ()
{
    int * D.2048;
    int * D.2049;
    int * p_a;
    int a[3];

    p_a = &a[0];
    *p_a = 10;
    D.2048 = p_a + 4;
    *D.2048 = 20;
    D.2049 = p_a + 8;
    *D.2049 = 30;
}
```



GIMPLE: Pointer to Array

test.c

```
int main()
{
    int *p_a, a[3];

    p_a = &a[0];

    *p_a = 10;
    *(p_a+1) = 20;
    *(p_a+2) = 30;
}
```

test.c.004t.gimple

```
main ()
{
    int * D.2048;
    int * D.2049;
    int * p_a;
    int a[3];

    p_a = &a[0];
    *p_a = 10;
    D.2048 = p_a + 4;
    *D.2048 = 20;
    D.2049 = p_a + 8;
    *D.2049 = 30;
}
```



GIMPLE: Pointer to Array

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int main()
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    int *p_a, a[3];

    p_a = &a[0];

    *p_a = 10;
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    *(p_a+2) = 30;
}
```

test.c.004t.gimple

```
main ()
{
    int * D.2048;
    int * D.2049;
    int * p_a;
    int a[3];

    p_a = &a[0];
    *p_a = 10;
    D.2048 = p_a + 4;
    *D.2048 = 20;
    D.2049 = p_a + 8;
    *D.2049 = 30;
}
```



GIMPLE: Pointer to Array

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int main()
{
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test.c.004t.gimple

```
main ()
{
    int * D.2048;
    int * D.2049;
    int * p_a;
    int a[3];

    p_a = &a[0];
    *p_a = 10;
    D.2048 = p_a + 4;
    *D.2048 = 20;
    D.2049 = p_a + 8;
    *D.2049 = 30;
}
```



GIMPLE: Translation of Conditional Statements

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```



GIMPLE: Translation of Conditional Statements

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```



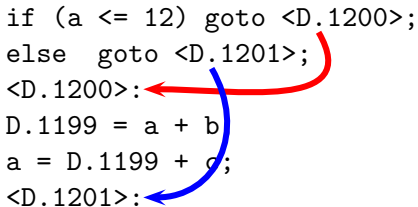
GIMPLE: Translation of Conditional Statements

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b
a = D.1199 + c;
<D.1201>:
```



GIMPLE: Translation of Loops

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
goto <D.1197>;
<D.1196>:
a = a + 1;
<D.1197>:
if (a <= 7) goto <D.1196>;
else goto <D.1198>;
<D.1198>:
```




GIMPLE: Translation of Loops

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
goto <D.1197>;
<D.1196>:
a = a + 1;
<D.1197>:
if (a <= 7) goto <D.1196>;
else goto <D.1198>;
<D.1198>:
```



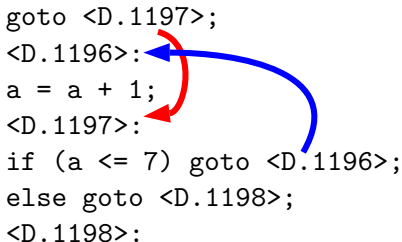
GIMPLE: Translation of Loops

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
goto <D.1197>;
<D.1196>:
a = a + 1;
<D.1197>:
if (a <= 7) goto <D.1196>;
else goto <D.1198>;
<D.1198>:
```



GIMPLE: Translation of Loops

test.c

```
int main()
{
    int a=2, b=3, c=4;
    while (a<=7)
    {
        a = a+1;
    }
    if (a<=12)
        a = a+b+c;
}
```

test.c.004t.gimple

```
goto <D.1197>;
<D.1196>:
a = a + 1;
<D.1197>:
if (a <= 7) goto <D.1196>;
else goto <D.1198>;
<D.1198>:
```



Control Flow Graph: Textual View

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```

test.c.013t.cfg

```
<bb 5>:
  if (a <= 12)
    goto <bb 6>;
  else
    goto <bb 7>;

<bb 6>:
  D.1199 = a + b;
  a = D.1199 + c;

<bb 7>:
  return;
```



Control Flow Graph: Textual View

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```

test.c.013t.cfg

```
<bb 5>:
  if (a <= 12)
    goto <bb 6>;
  else
    goto <bb 7>;

<bb 6>:
  D.1199 = a + b;
  a = D.1199 + c;

<bb 7>:
  return;
```



Control Flow Graph: Textual View

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```

test.c.013t.cfg

```
<bb 5>:
  if (a <= 12)
    goto <bb 6>;
  else
    goto <bb 7>;

<bb 6>:
  D.1199 = a + b;
  a = D.1199 + c;

<bb 7>:
  return;
```



Control Flow Graph: Textual View

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```

test.c.013t.cfg

```
<bb 5>:
  if (a <= 12)
    goto <bb 6>;
  else
    goto <bb 7>;

<bb 6>:
  D.1199 = a + b;
  a = D.1199 + c;

<bb 7>:
  return;
```



Control Flow Graph: Textual View

test.c.004t.gimple

```
if (a <= 12) goto <D.1200>;
else goto <D.1201>;
<D.1200>:
D.1199 = a + b;
a = D.1199 + c;
<D.1201>:
```

test.c.013t.cfg

```
<bb 5>:
  if (a <= 12)
    goto <bb 6>;
  else
    goto <bb 7>;

<bb 6>:
  D.1199 = a + b;
  a = D.1199 + c;

<bb 7>:
  return;
```



OC

A. Koch

Compile-Fluß

Aufbau

GIMPLE IR

RTL IR

Optimierung

Zwischendarstellung RTL

Part 3

Examining RTL Dumps

RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: test.c.144r.expand

```
(insn 12 11 13 4 (parallel [  
  ( set (mem/c/i:SI  
        (plus:SI  
          (reg/f:SI 54 virtual-stack-vars)  
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))  
  (plus:SI  
    (mem/c/i:SI  
      (plus:SI  
        (reg/f:SI 54 virtual-stack-vars)  
        (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))  
    (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t.c:24 -1 (nil))
```

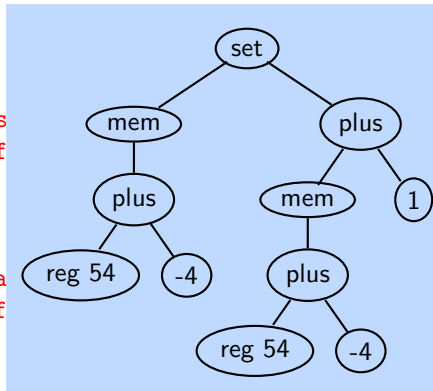


RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: test.c.144r.expand

```
(insn 12 11 13 4 (parallel [
  ( set (mem/c/i:SI
    (plus:SI
      (reg/f:SI 54 virtual-s
      (const_int -4 [0xffff
    (plus:SI
      (mem/c/i:SI
      (plus:SI
        (reg/f:SI 54 virtua
        (const_int -4 [0xff
      (const_int 1 [0x1])))
    (clobber (reg:CC 17 flags))
  ]) t.c:24 -1 (nil))
```

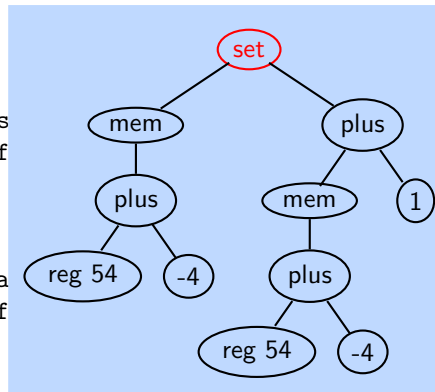


RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: test.c.144r.expand

```
(insn 12 11 13 4 (parallel [
  ( set (mem/c/i:SI
    (plus:SI
      (reg/f:SI 54 virtual-s
      (const_int -4 [0xffff
    (plus:SI
      (mem/c/i:SI
      (plus:SI
        (reg/f:SI 54 virtua
        (const_int -4 [0xff
      (const_int 1 [0x1])))
    (clobber (reg:CC 17 flags))
  ]) t.c:24 -1 (nil))
```



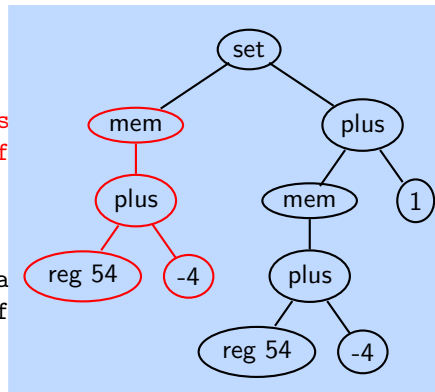
RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: test.c.144r.expand

```
(insn 12 11 13 4 (parallel [
  ( set (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-s
            (const_int -4 [0xffff
              (plus:SI
                (mem/c/i:SI
                  (plus:SI
                    (reg/f:SI 54 virtua
                      (const_int -4 [0xff
                        (const_int 1 [0x1])))
                    (clobber (reg:CC 17 flags))
                  ]) t.c:24 -1 (nil))
```

a is a local variable
allocated on stack



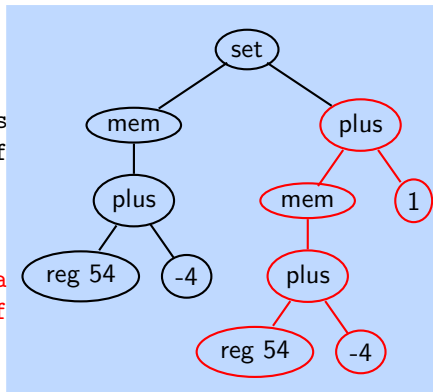
RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: test.c.144r.expand

```
(insn 12 11 13 4 (parallel [
  ( set (mem/c/i:SI
    (plus:SI
      (reg/f:SI 54 virtual-s
        (const_int -4 [0xffff
          (plus:SI
            (mem/c/i:SI
              (plus:SI
                (reg/f:SI 54 virtua
                  (const_int -4 [0xff
                    (const_int 1 [0x1])))
                (clobber (reg:CC 17 flags))
              ]) t.c:24 -1 (nil))
```

a is a local variable
allocated on stack



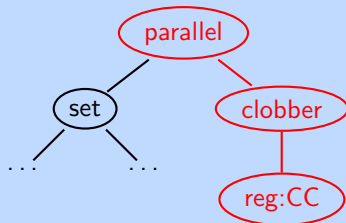
RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: `test.c.144r.expand`

```
(insn 12 11 13 4 (parallel [
  ( set (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-s
            (const_int -4 [0xffff
              (plus:SI
                (mem/c/i:SI
                  (plus:SI
                    (reg/f:SI 54 virtua
                      (const_int -4 [0xff
                        (const_int 1 [0x1])))
                  (clobber (reg:CC 17 flags))
                ]) t.c:24 -1 (nil))
```

side-effect of plus may
modify condition code register
non-deterministically



RTL for i386: Arithmetic Operations (1)

Translation of $a = a + 1$

Dump file: test.c.144r.e Output with slim suffix

```
(insn 12 11 13 4 (parallel
  ( set (mem/c/i:SI
    (plus:SI
      (reg/f:SI 54 virtual-stack-vars)
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])
    (plus:SI
      (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-stack-vars)
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])
          (const_int 1 [0x1])))
      (clobber (reg:CC 17 flags))
    ]) t.c:24 -1 (nil))
  { [r54:SI-0x4]=[r54:SI-0x4]+0x1;
    clobber flags:CC;
  }
```



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
  (plus:SI  
    (mem/c/i:SI  
      (plus:SI  
        (reg/f:SI 54 virtual-stack-vars)  
        (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
    (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t.c:24 -1 (nil))
```



Additional Information in RTL

```

(insn 12 11 13 4 (parallel [
  (set (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-stack-vars)
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))
    (plus:SI
      (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-stack-vars)
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))
      (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:24 -1 (nil))

```

Current Instruction



Additional Information in RTL

```

(insn 12 11 13 4 (parallel [
  (set (mem/c/i:SI
    (plus:SI
      (reg/f:SI 54 virtual-stack-vars)
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))
    (plus:SI
      (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-stack-vars)
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))
        (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:24 -1 (nil))

```

Previous Instruction



Additional Information in RTL

```

(insn 12 11 13 4 (parallel [
  (set (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-stack-vars)
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))
    (plus:SI
      (mem/c/i:SI
        (plus:SI
          (reg/f:SI 54 virtual-stack-vars)
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32]))
      (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:24 -1 (nil))

```


Next Instruction



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
  (plus:SI  
    (mem/c/i:SI  
      (plus:SI  
        (reg/f:SI 54 virtual-stack-vars)  
        (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
    (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t.c:24 -1 (nil))
```

Basic Block



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
  (plus:SI  
    (mem/c/i:SI  
      (plus:SI  
        (reg/f:SI 54 virtual-stack-vars)  
        (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
    (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t.c:24 -1 (nil))
```

File name: Line number



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
    (plus:SI  
      (mem/c/i:SI  
        (plus:SI  
          (reg/f:SI 54 virtual-stack-vars)  
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
        (const_int 1 [0x1])))  
    (clobber (reg:CC 17 flags))  
  ]) t.c:24 -1 (nil))
```

memory reference
that does not trap



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
    (plus:SI  
      (mem/c/i:SI  
        (plus:SI  
          (reg/f:SI 54 virtual-stack-vars)  
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
        (const_int 1 [0x1])))  
    (clobber (reg:CC 17 flags))  
  ]) t.c:24 -1 (nil))
```

scalar that is not a part of an aggregate



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
  (plus:SI  
    (mem/c/i:SI  
      (plus:SI  
        (reg/f:SI 54 virtual-stack-vars)  
        (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
      (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t.c:24 -1 (nil))
```

register that holds a pointer



Additional Information in RTL

```
(insn 12 11 13 4 (parallel [  
  (set (mem/c/i:SI  
    (plus:SI  
      (reg/f:SI 54 virtual-stack-vars)  
      (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
    (plus:SI  
      (mem/c/i:SI  
        (plus:SI  
          (reg/f:SI 54 virtual-stack-vars)  
          (const_int -4 [0xffffffffc])) [0 a+0 S4 A32])  
        (const_int 1 [0x1])))  
    (clobber (reg:CC 17 flags))  
  ]) t.c:24 -1 (nil))
```

single integer



RTL for i386: Arithmetic Operations (2)

Translation of $a = a + 1$ when a is a global variable

Dump file: test.c.144r.expand

```
(insn 11 10 12 4 (set
  (reg:SI 64 [ a.0 ])
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+0 S4 A32]))) t.c:26 -1 (nil))
```

```
(insn 12 11 13 4 (parallel [
  (set (reg:SI 63 [ a.1 ])
    (plus:SI (reg:SI 64 [ a.0 ])
      (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:26 -1 (nil))
```

```
(insn 13 12 14 4 (set
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+0 S4 A32])
  (reg:SI 63 [ a.1 ])) t.c:26 -1 (nil))
```



RTL for i386: Arithmetic Operations (2)

Translation of $a = a + 1$ when a is a global variable

Dump file: test.c.144r.expand

```
(insn 11 10 12 4 (set
  (reg:SI 64 [ a.0 ])
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+
  )

(insn 12 11 13 4 (parallel [
  (set (reg:SI 63 [ a.1 ])
    (plus:SI (reg:SI 64 [ a.0 ])
      (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:26 -1 (nil))

(insn 13 12 14 4 (set
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+0 S4 A32])
  (reg:SI 63 [ a.1 ])) t.c:26 -1 (nil))
```

Load a into reg64



RTL for i386: Arithmetic Operations (2)

Translation of $a = a + 1$ when a is a global variable

Dump file: test.c.144r.expand

```
(insn 11 10 12 4 (set
  (reg:SI 64 [ a.0 ])
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+
    )
  (plus:SI (reg:SI 64 [ a.0 ])
    (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:26 -1 (nil))

(insn 12 11 13 4 (parallel [
  (set (reg:SI 63 [ a.1 ])
    (plus:SI (reg:SI 64 [ a.0 ])
      (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:26 -1 (nil))

(insn 13 12 14 4 (set
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+0 S4 A32])
  (reg:SI 63 [ a.1 ])) t.c:26 -1 (nil))
```

Load a into reg64

reg63 = reg64 + 1



RTL for i386: Arithmetic Operations (2)

Translation of $a = a + 1$ when a is a global variable

Dump file: test.c.144r.expand

```
(insn 11 10 12 4 (set
  (reg:SI 64 [ a.0 ])
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+
    )
  )

(insn 12 11 13 4 (parallel [
  (set (reg:SI 63 [ a.1 ])
    (plus:SI (reg:SI 64 [ a.0 ])
      (const_int 1 [0x1])))
  (clobber (reg:CC 17 flags))
]) t.c:26 -1 (nil))

(insn 13 12 14 4 (set
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+0 S4 A32])
  (reg:SI 63 [ a.1 ])) t.c:26 -1 (nil))
```

Load a into reg64
 reg63 = reg64 + 1
 store reg63 into a



RTL for i386: Arithmetic Operations (2)

Translation of $a = a + 1$ when a is a global variable

Dump file: test.c.144r.expand

```
(insn 11 10 12 4 (set
  (reg:SI 64 [ a.0 ])
  (mem/c/i:SI (symbol_ref:SI ("a")
    <var_decl 0xb7d8d000 a>) [0 a+
    )
  )
  (insn 12 11 13 4 (parallel [
    (set (reg:SI 63 [ a.1 ])
      (plus:SI (reg:SI 64 [ a.0 ])
        (const_int 1 [0x1])))
    (clobber (reg:CC 17 flags))
  ]) t.c:26 -1 (nil))
  (insn 13 12 14 4 (set
    (mem/c/i:SI (symbol_ref:SI ("a")
      <var_decl 0xb7d8d000 a>) [0 a+0 S4 A32])
    (reg:SI 63 [ a.1 ])) t.c:26 -1 (nil))
```

Load a into reg64
 reg63 = reg64 + 1
 store reg63 into a

Output with slim suffix
 r64:SI=['a']
 {r63:SI=r64:SI+0x1;
 clobber flags:CC;
 }
 ['a']=r63:SI



RTL for i386: Arithmetic Operations (3)

Translation of $a = a + 1$ when a is a formal parameter

Dump file: test.c.144r.expand

```
(insn 10 9 11 4 (parallel [  
  (set  
    (mem/c/i:SI  
      (reg/f:SI 53 virtual-incoming-args) [0 a+0 S4 A32]))  
    (plus:SI  
      (mem/c/i:SI  
        (reg/f:SI 53 virtual-incoming-args) [0 a+0 S4 A32]))  
      (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t1.c:25 -1 (nil))
```



RTL for i386: Arithmetic Operations (3)

Translation of $a = a + 1$ when a is a formal parameter

Dump file: test.c.144r.expand

```
(insn 10 9 11 4 (parallel [
  (set
    (mem/c/i:SI
      (reg/f:SI 53 virtual-incoming-
        (plus:SI
          (mem/c/i:SI
            (reg/f:SI 53 virtual-incoming-
              (const_int 1 [0x1])))
          (clobber (reg:CC 17 flags))
        ])) t1.c:25 -1 (nil))
```

Access through argument pointer register instead of frame pointer register



RTL for i386: Arithmetic Operations (3)

Translation of $a = a + 1$ when a is a formal parameter

Dump file: test.c.144r.expand

```
(insn 10 9 11 4 (parallel [
  (set
    (mem/c/i:SI
      (reg/f:SI 53 virtual-incoming-
        (plus:SI
          (mem/c/i:SI
            (reg/f:SI 53 virtual-incoming-
              (const_int 1 [0x1])))
          (clobber (reg:CC 17 flags))
        ])) t1.c:25 -1 (nil))
```

Access through argument
pointer register instead of
frame pointer register

No offset required?



RTL for i386: Arithmetic Operations (3)

Translation of $a = a + 1$ when a is a formal parameter

Dump file: test.c.144r.expand

```
(insn 10 9 11 4 (parallel [
  (set
    (mem/c/i:SI
      (reg/f:SI 53 virtual-incoming-
        (plus:SI
          (mem/c/i:SI
            (reg/f:SI 53 virtual-incoming-
              (const_int 1 [0x1])))
          (clobber (reg:CC 17 flags))
        ])) t1.c:25 -1 (nil))
```

Access through argument pointer register instead of frame pointer register

No offset required?

Output with slim suffix

```
{[r53:SI]=[r53:SI]+0x1;
clobber flags:CC;
}
```



RTL for i386: Arithmetic Operation (4)

Translation of $a = a + 1$ when **a** is the second formal parameter

Dump file: `test.c.144r.expand`

```
(insn 10 9 11 4 (parallel [  
  (set  
    (mem/c/i:SI  
      (plus:SI  
        (reg/f:SI 53 virtual-incoming-args)  
        (const_int 4 [0x4]))) [0 a+0 S4 A32])  
    (plus:SI  
      (mem/c/i:SI  
        (plus:SI  
          (reg/f:SI 53 virtual-incoming-args)  
          (const_int 4 [0x4]))) [0 a+0 S4 A32])  
      (const_int 1 [0x1])))  
  (clobber (reg:CC 17 flags))  
) t1.c:25 -1 (nil))
```



RTL for i386: Arithmetic Operation (4)

Translation of $a = a + 1$ when a is the second formal parameter

Dump file: `test.c.144r.expand`

```
(insn 10 9 11 4 (parallel [
  (set
    (mem/c/i:SI
      (plus:SI
        (reg/f:SI 53 virtual-
          (const_int 4 [0x4]))
      (plus:SI
        (mem/c/i:SI
          (plus:SI
            (reg/f:SI 53 virtu
              (const_int 4 [0x4]
            (const_int 1 [0x1])))
        (clobber (reg:CC 17 flags))
    ]) t1.c:25 -1 (nil))
```

Offset 4 added to the argument pointer register



RTL for i386: Arithmetic Operation (4)

Translation of $a = a + 1$ when a is the second formal parameter

Dump file: `test.c.144r.expand`

```
(insn 10 9 11 4 (parallel [
  (set
    (mem/c/i:SI
      (plus:SI
        (reg/f:SI 53 virtual-
          (const_int 4 [0x4]))
        (plus:SI
          (mem/c/i:SI
            (plus:SI
              (reg/f:SI 53 virtu
                (const_int 4 [0x4]
              (const_int 1 [0x1]))))
          (clobber (reg:CC 17 flags))
    ]) t1.c:25 -1 (nil))
```

Offset 4 added to the argument pointer register

When a is the first parameter, its offset is 0!



RTL for i386: Arithmetic Operation (4)

Translation of $a = a + 1$ when a is the second formal parameter

Dump file: `test.c.144r.expand`

```
(insn 10 9 11 4 (parallel [
  (set
    (mem/c/i:SI
      (plus:SI
        (reg/f:SI 53 virtual-
          (const_int 4 [0x4])))
      (plus:SI
        (mem/c/i:SI
          (plus:SI
            (reg/f:SI 53 virtu
              (const_int 4 [0x4]
                (const_int 1 [0x1]))))
          (clobber (reg:CC 17 flags))
        ]) t1.c:25 -1 (nil))
```

Offset 4 added to the argument pointer register

When a is the first parameter, its offset is 0!

Output with slim suffix

```
{ [r53:SI+0x4]=[r53:SI+0x4]+0x1;
  clobber flags:CC;
}
```



RTL for spim: Arithmetic Operations

Translation of $a = a + 1$ when a is a local variable

Dump file: `test.c.144r.expand`

```
r39=stack($fp - 4)
r40=r39+1
stack($fp - 4)=r40
```

```
(insn 7 6 8 4 (set (reg:SI 39)
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])) -1 (nil))
(insn 8 7 9 4 test.c:6 (set (reg:SI 40)
  (plus:SI (reg:SI 39)
    (const_int 1 [...]))) -1 (nil))
(insn 9 8 10 4 test.c:6 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])
  (reg:SI 40)) test.c:6 -1 (nil))
```

In spim, a variable is loaded into register to perform any instruction, hence three instructions are generated



RTL for spim: Arithmetic Operations

Translation of $a = a + 1$ when a is a local variable

Dump file: test.c.144r.expand

```
r39=stack($fp - 4)
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stack($fp - 4)=r40
```

```
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  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])) -1 (nil))
(insn 8 7 9 4 test.c:6 (set (reg:SI 40)
  (plus:SI (reg:SI 39)
    (const_int 1 [...]))) -1 (nil))
(insn 9 8 10 4 test.c:6 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])
  (reg:SI 40)) test.c:6 -1 (nil))
```

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RTL for spim: Arithmetic Operations

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stack($fp - 4)=r40
```

```
(insn 7 6 8 4 (set (reg:SI 39)
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])) -1 (nil))
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  (plus:SI (reg:SI 39)
    (const_int 1 [...]))) -1 (nil))
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    (const_int -4 [...])) [...])
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r40=r39+1
stack($fp - 4)=r40
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(insn 7 6 8 4 (set (reg:SI 39)
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])) -1 (nil))
(insn 8 7 9 4 test.c:6 (set (reg:SI 40)
  (plus:SI (reg:SI 39)
    (const_int 1 [...]))) -1 (nil))
(insn 9 8 10 4 test.c:6 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 33 virtual-stack-vars)
    (const_int -4 [...])) [...])
  (reg:SI 40)) test.c:6 -1 (nil))
```

In spim, a variable is loaded into register to perform any instruction, hence three instructions are generated



RTL for i386: Control Flow

What does this represent?

```
(jump_insn 15 14 16 4 (set (pc)
  (if_then_else (lt (reg:CCGC 17 flags)
    (const_int 0 [0x0]))
    (label_ref 12)
    (pc))) p1.c:6 -1 (nil)
(nil)
-> 12)
```



RTL for i386: Control Flow

What does this represent?

```
(jump_insn 15 14 16 4 (set (pc)
  (if_then_else (lt (reg:CCGC 17 flags)
    (const_int 0 [0x0]))
    (label_ref 12)
    (pc))) p1.c:6 -1 (nil)
(nil)
-> 12)
```

$pc = r17 < 0 ? \text{label}(12) : pc$



RTL for i386: Control Flow

Translation of `if (a > b) { /* something */ }`

Dump file: `test.c.144r.expand`

```
(insn 8 7 9 (set (reg:SI 61)
  (mem/c/i:SI (plus:SI (reg/f:SI 54 virtual-stack-vars)
    (const_int -8 [0xffffffff]))) [0 a+0 S4 A32])) test.c:7 -1
(insn 9 8 10 (set (reg:CCGC 17 flags)
  (compare:CCGC (reg:SI 61)
    (mem/c/i:SI (plus:SI (reg/f:SI 54 virtual-stack-vars)
      (const_int -4 [0xfffffff8]))) [0 b+0 S4 A32]))) test.c:7
(jump_insn 10 9 0 (set (pc)
  (if_then_else (le (reg:CCGC 17 flags)
    (const_int 0 [0x0]))
    (label_ref 13)
    (pc))) test.c:7 -1 (nil)
-> 13)
```



RTL for i386: Control Flow

Translation of `if (a > b) { /* something */ }`

Dump file: `test.c.144r.expand`

```
(insn 8 7 9 (set (reg:SI 61)
  (mem/c/i:SI (plus:SI (reg/f:SI 54 virtual-stack-vars)
    (const_int -8 [0xffffffff]))) [0 a+0 S4 A32])) test.c:7 -1
(insn 9 8 10 (set (reg:CCGC 17 flags)
  (compare:CCGC (reg:SI 61)
    (mem/c/i:SI (plus:SI (reg/f:SI 54 virtual-stack-vars)
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RTL for i386: Control Flow

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-> 13)
```



Observing Register Allocation for i386

test.c

```
int main()
{
    int a=2, b=3;
    if(a<=12)
        a = a * b;
}
```

test.c.188r.asmcons

(observable dump before register allocation)

```
(insn 10 9 11 3 (set (reg:SI 59)
  (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
    (const_int -4 [0xffffffffc]))) [0 a+0 S4 A32

(insn 11 10 12 3 (parallel [
  (set (reg:SI 60)
    (mult:SI (reg:SI 59)
      (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
        (const_int -8 [0xffffffff8]))) [0 b+0
  (clobber (reg:CC 17 flags))
]) 262 *mulsi3_1 test.c:5 (nil))

(insn 12 11 22 3 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
    (const_int -4 [0xffffffffc]))) [0 a+0 S4
  (reg:SI 60)) 44 *movsi_internal test.c:
```



Observing Register Allocation for i386

```
test.c
```

```
int main()
{
  int a=2, b=3;
  if(a<=12)
    a = a * b;
}
```

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test.c.188r.asmcons
```

```
(observable dump before register allocation)
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```
(insn 10 9 11 3 (set (reg:SI 59)
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    (const_int -4 [0xffffffffc]))) [0 a+0 S4 A32
```

```
(insn 11 10 12 3 (parallel [
  (set (reg:SI 60)
    (mult:SI (reg:SI 59)
      (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
        (const_int -8 [0xffffffff8]))) [0 b+0
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]) 262 *mulsi3_1 test.c:5 (nil))
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```
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```



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    (const_int -4 [0xffffffffc]))) [0 a+0 S4 A32

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  (set (reg:SI 60)
    (mult:SI (reg:SI 59)
      (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
        (const_int -8 [0xffffffff8]))) [0 b+0
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```



Observing Register Allocation for i386

```
test.c
```

```
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{
    int a=2, b=3;
    if(a<=12)
        a = a * b;
}
```

```
test.c.188r.asmcons
```

```
(observable dump before register allocation)
```

```
(insn 10 9 11 3 (set (reg:SI 59)
  (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
    (const_int -4 [0xffffffffc]))) [0 a+0 S4 A32
```

```
(insn 11 10 12 3 (parallel [
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]) 262 *mulsi3_1 test.c:5 (nil))
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```
(insn 12 11 22 3 (set
  (mem/c/i:SI (plus:SI (reg/f:SI 20 frame)
    (const_int -4 [0xffffffffc]))) [0 a+0 S4
  (reg:SI 60)) 44 *movsi_internal test.c:
```



Observing Register Allocation for i386

test.c.188r.asmcons

```
(set (reg:SI 59) (mem/c/i:SI
  (plus:SI
    (reg/f:SI 20 frame)
    (const_int -4))))
```

```
(set (reg:SI 60)
  (mult:SI
    (reg:SI 59)
    (mem/c/i:SI
      (plus:SI
        (reg/f:SI 20 frame)
        (const_int -8)) )))
```

```
(set (mem/c/i:SI (plus:SI
  (reg/f:SI 20 frame)
  (const_int -4)))
  (reg:SI 60))
```

test.c.188r.ira

```
(set (reg:SI 0 ax [59]) (mem/c/i:SI
  (plus:SI
    (reg/f:SI 6 bp)
    (const_int -4))))
```

```
(set (reg:SI 0 ax [60])
  (mult:SI
    (reg:SI 0 ax [59])
    (mem/c/i:SI
      (plus:SI
        (reg/f:SI 6 bp)
        (const_int -8)) )))
```

```
(set (mem/c/i:SI (plus:SI
  (reg/f:SI 6 bp)
  (const_int -4)))
  (reg:SI 0 ax [60]))
```



OC

A. Koch

Compile-Fluß

Aufbau

GIMPLE IR

RTL IR

Optimierung

Maschinenunabhängige Optimierung

Outline

- Example 1
 - ▶ Constant Propagation
 - ▶ Copy Propagation
 - ▶ Dead Code Elimination
 - ▶ Loop unrolling
- Example 2
 - ▶ Partial Redundancy Elimination
 - ▶ Copy Propagation
 - ▶ Dead Code Elimination



Part 1

First Example Program

Example Program 1

```
int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

- What does this program return?



Example Program 1

```
int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

- What does this program return?
- 12



Example Program 1

```
int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

- What does this program return?
- 12
- We use this program to illustrate various shades of the following optimizations:
Constant propagation, Copy propagation, Loop unrolling, Dead code elimination



Compilation Command

```
$gcc -fdump-tree-all -O2 ccp.c
```



Example Program 1

Program ccp.c

```
int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

Control flow graph



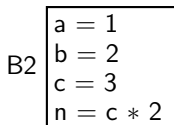
Example Program 1

Program ccp.c

```
int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

Control flow graph



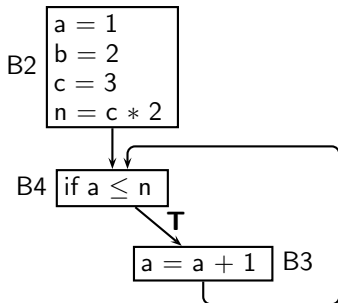
Example Program 1

Program ccp.c

```
int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

Control flow graph



Example Program 1

Program ccp.c

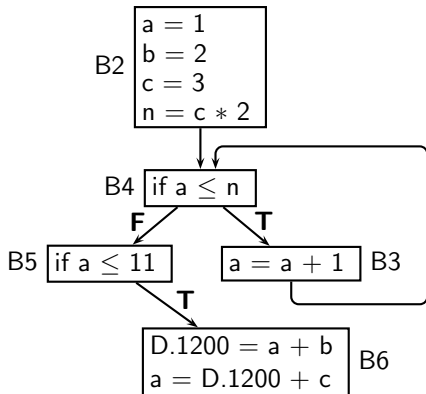
```

int main()
{ int a, b, c, n;

  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}

```

Control flow graph



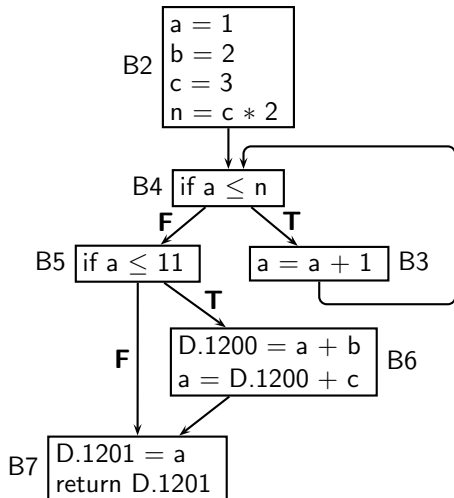
Example Program 1

Program ccp.c

```
int main()
{ int a, b, c, n;

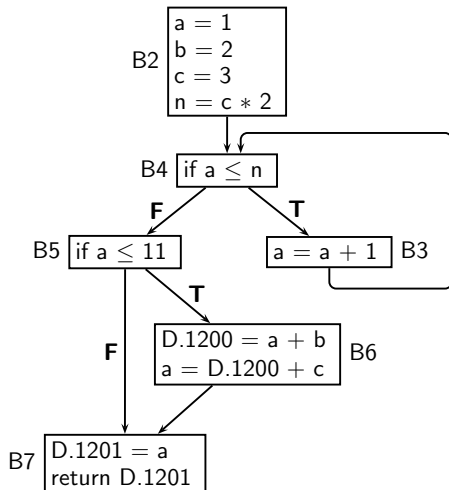
  a = 1;
  b = 2;
  c = 3;
  n = c*2;
  while (a <= n)
  {
    a = a+1;
  }
  if (a < 12)
    a = a+b+c;
  return a;
}
```

Control flow graph



Control Flow Graph: Pictorial and Textual View

Control flow graph

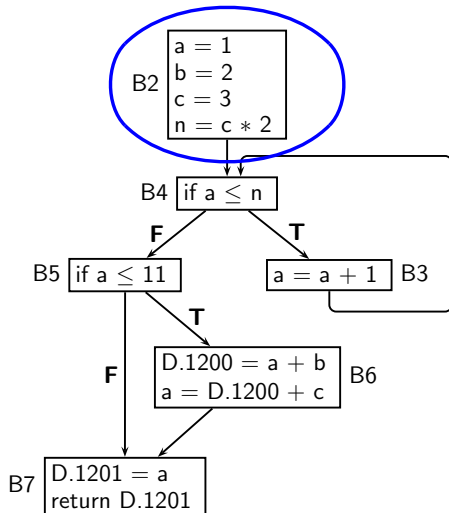


Dump file `ccp.c.013t.cfg`



Control Flow Graph: Pictorial and Textual View

Control flow graph



Dump file ccp.c.013t.cfg

```

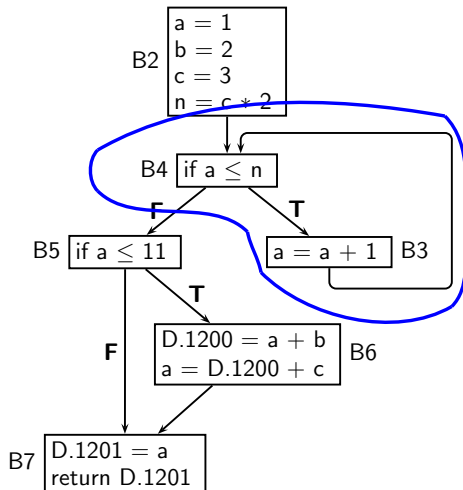
<bb 2>:
a = 1;
b = 2;
c = 3;
n = c * 2;
goto <bb 4>;

```



Control Flow Graph: Pictorial and Textual View

Control flow graph



Dump file ccp.c.013t.cfg

```

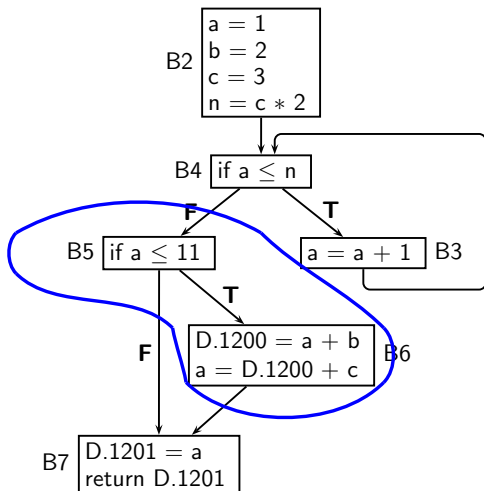
<bb 3>:
a = a + 1;

<bb 4>:
if (a <= n)
    goto <bb 3>;
else
    goto <bb 5>;
  
```



Control Flow Graph: Pictorial and Textual View

Control flow graph

Dump file `ccp.c.013t.cfg`

```

<bb 5>:
if (a <= 11)
    goto <bb 6>;
else
    goto <bb 7>;

```

```

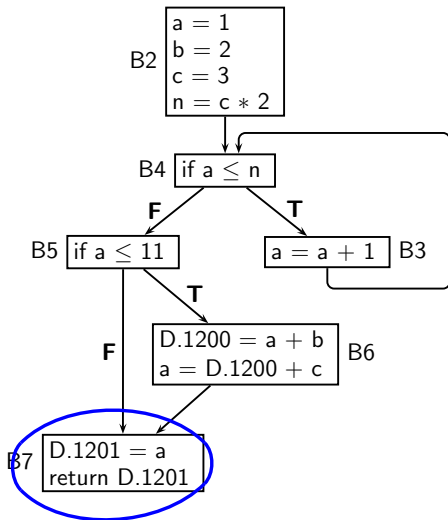
<bb 6>:
D.1200 = a + b;
a = D.1200 + c;

```



Control Flow Graph: Pictorial and Textual View

Control flow graph



Dump file `ccp.c.013t.cfg`

<bb 7>:

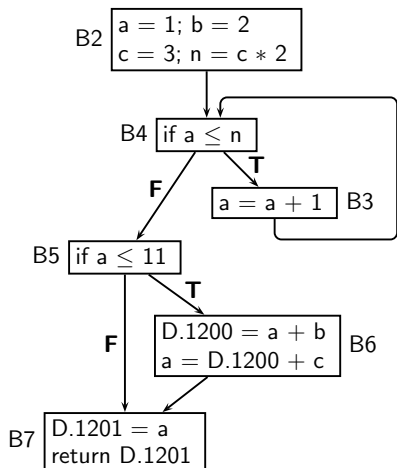
```
D.1201 = a;
return D.1201;
```



Single Static Assignment (SSA) Form

Control flow graph

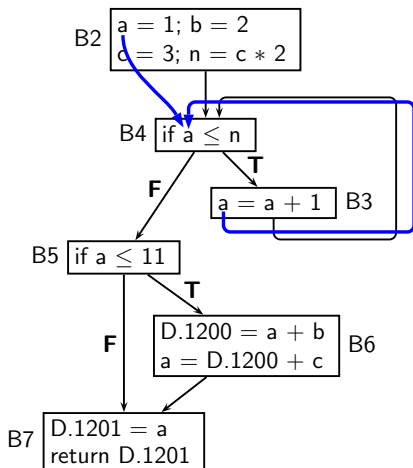
SSA Form



Single Static Assignment (SSA) Form

Control flow graph

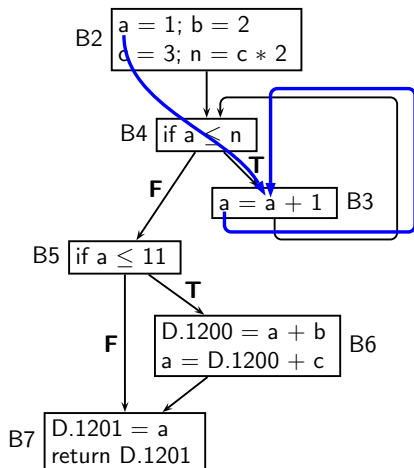
SSA Form



Single Static Assignment (SSA) Form

Control flow graph

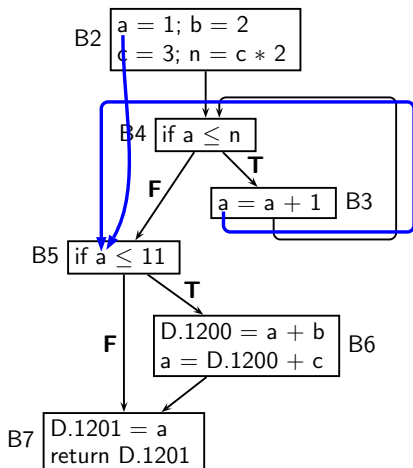
SSA Form



Single Static Assignment (SSA) Form

Control flow graph

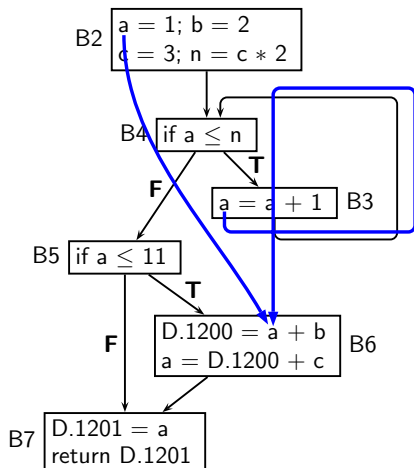
SSA Form



Single Static Assignment (SSA) Form

Control flow graph

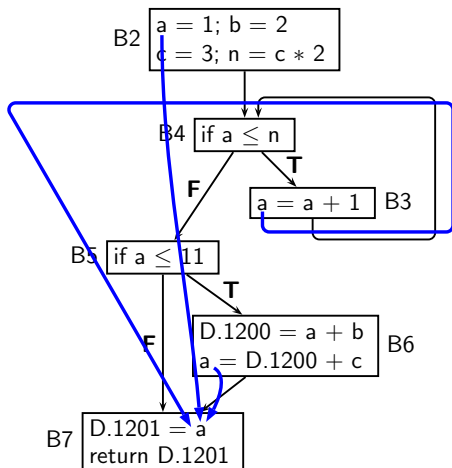
SSA Form



Single Static Assignment (SSA) Form

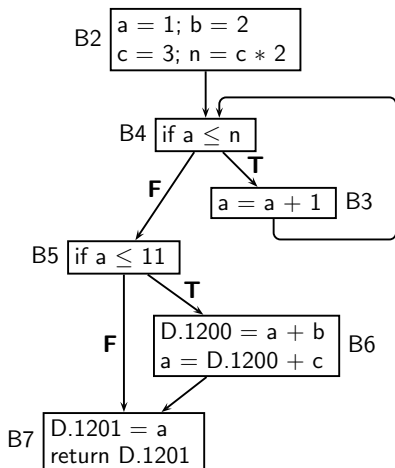
Control flow graph

SSA Form

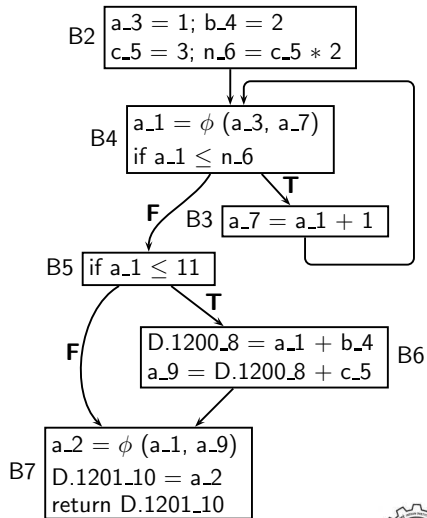


Single Static Assignment (SSA) Form

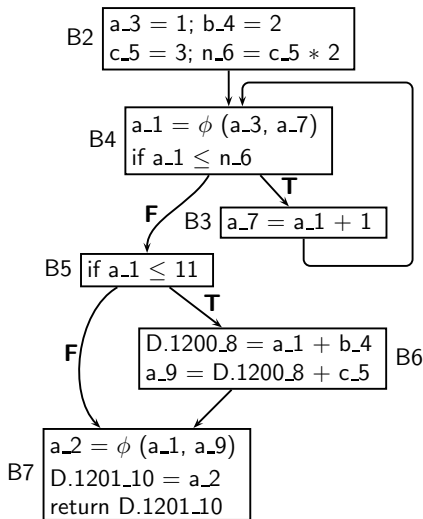
Control flow graph



SSA Form



Properties of SSA Form



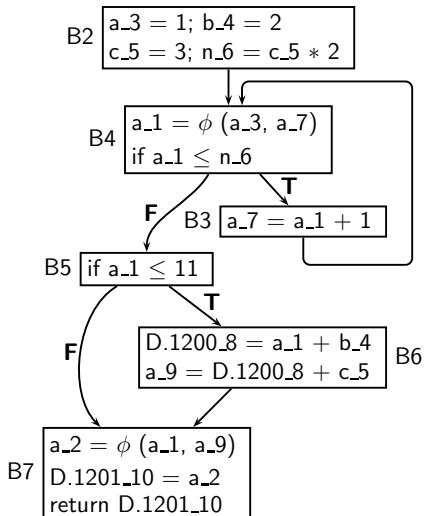
- A ϕ function is a multiplexer or a selection function
- Every use of a variable corresponds to a unique definition of the variable
- For every use, the definition is guaranteed to appear on every path leading to the use

SSA construction algorithm is expected to insert as few ϕ functions as possible to ensure the above properties



SSA Form: Pictorial and Textual View

CFG in SSA form

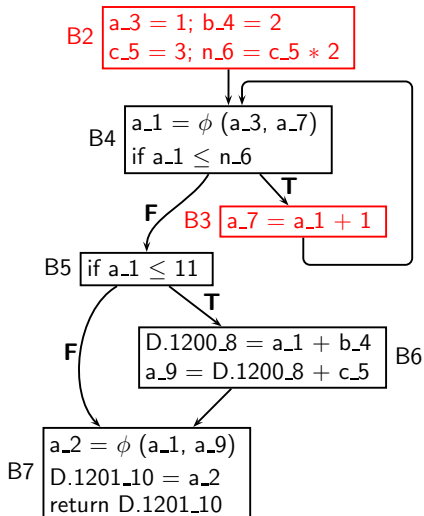


Dump file [ccp.c.017t.ssa](#)



SSA Form: Pictorial and Textual View

CFG in SSA form



Dump file `ccp.c.017t.ssa`

<bb 2>:

`a_3 = 1;`

`b_4 = 2;`

`c_5 = 3;`

`n_6 = c_5 * 2;`

`goto <bb 4>;`

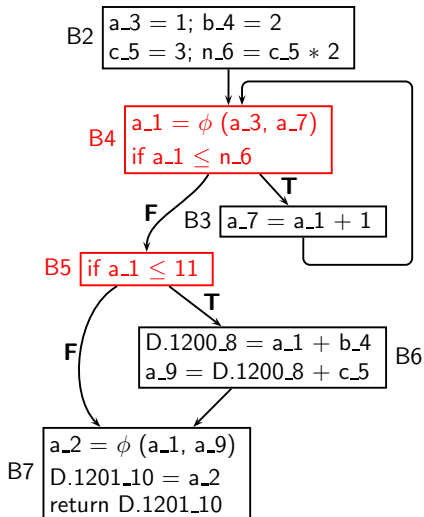
<bb 3>:

`a_7 = a_1 + 1;`



SSA Form: Pictorial and Textual View

CFG in SSA form



Dump file `ccp.c.017t.ssa`

<bb 4>:

```
# a_1 = PHI <a_3(2), a_7(3)>
if (a_1 <= n_6)
  goto <bb 3>;
else
  goto <bb 5>;
```

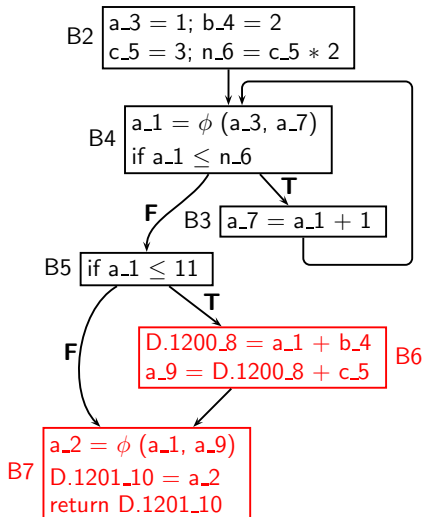
<bb 5>:

```
if (a_1 <= 11)
  goto <bb 6>;
else
  goto <bb 7>;
```



SSA Form: Pictorial and Textual View

CFG in SSA form



Dump file `ccp.c.017t.ssa`

<bb 6>:

```
D.1200_8 = a_1 + b_4;
a_9 = D.1200_8 + c_5;
```

<bb 7>:

```
# a_2 = PHI <a_1(5), a_9(6)>
D.1201_10 = a_2;
return D.1201_10;
```



A Comparison of CFG and SSA Dumps

Dump file ccp.c.013t.cfg

Dump file ccp.c.017t.ssa



A Comparison of CFG and SSA Dumps

Dump file ccp.c.013t.cfg

```
<bb 2>:  
  a = 1;  
  b = 2;  
  c = 3;  
  n = c * 2;  
  goto <bb 4>;
```

```
<bb 3>:  
  a = a + 1;
```

Dump file ccp.c.017t.ssa

```
<bb 2>:  
  a_3 = 1;  
  b_4 = 2;  
  c_5 = 3;  
  n_6 = c_5 * 2;  
  goto <bb 4>;
```

```
<bb 3>:  
  a_7 = a_1 + 1;
```



A Comparison of CFG and SSA Dumps

Dump file ccp.c.013t.cfg

```
<bb 4>:
  if (a <= n)
    goto <bb 3>;
  else
    goto <bb 5>;

<bb 5>:
  if (a <= 11)
    goto <bb 6>;
  else
    goto <bb 7>;
```

Dump file ccp.c.017t.ssa

```
<bb 4>:
  # a_1 = PHI <a_3(2), a_7(3)>
  if (a_1 <= n_6)
    goto <bb 3>;
  else
    goto <bb 5>;

<bb 5>:
  if (a_1 <= 11)
    goto <bb 6>;
  else
    goto <bb 7>;
```



A Comparison of CFG and SSA Dumps

Dump file ccp.c.013t.cfg

```
<bb 6>:  
D.1200 = a + b;  
a = D.1200 + c;
```

```
<bb 7>:  
D.1201 = a;  
return D.1201;
```

Dump file ccp.c.017t.ssa

```
<bb 6>:  
  D.1200_8 = a_1 + b_4;  
  a_9 = D.1200_8 + c_5;
```

```
<bb 7>:  
  # a_2 = PHI <a_1(5), a_9(6)>  
  D.1201_10 = a_2;  
  return D.1201_10;
```



Copy Renaming

Input dump: ccp.c.017t.ssa

```
<bb 7>:  
# a_2 = PHI <a_1(5), a_9(6)>  
D.1201_10 = a_2;  
return D.1201_10;
```

Output dump: ccp.c.022t.copyrename1

```
<bb 7>:  
# a_2 = PHI <a_1(5), a_9(6)>  
a_10 = a_2;  
return a_10;
```



First Level Constant and Copy Propagation

Input dump: ccp.c.022t.copyrename1

```
<bb 2>:
  a_3 = 1;
  b_4 = 2;
  c_5 = 3;
  n_6 = c_5 * 2;
  goto <bb 4>;
```

```
<bb 3>:
  a_7 = a_1 + 1;
```

```
<bb 4>:
  # a_1 = PHI < a_3(2), a_7(3)>
  if (a_1 <= n_6)
    goto <bb 3>;
  else
    goto <bb 5>;
```

Output dump: ccp.c.023t.ccp1

```
<bb 2>:
  a_3 = 1;
  b_4 = 2;
  c_5 = 3;
  n_6 = 6;
  goto <bb 4>;
```

```
<bb 3>:
  a_7 = a_1 + 1;
```

```
<bb 4>:
  # a_1 = PHI < 1(2), a_7(3)>
  if (a_1 <= 6)
    goto <bb 3>;
  else
    goto <bb 5>;
```



First Level Constant and Copy Propagation

Input dump: ccp.c.022t.copyrename1

```
<bb 2>:  
  a_3 = 1;  
  b_4 = 2;  
  c_5 = 3;  
  n_6 = 6;  
  goto <bb 4>;
```

...

```
<bb 6>:  
  D.1200_8 = a_1 + b_4;  
  a_9 = D.1200_8 + c_5;
```

Output dump: ccp.c.023t.ccp1

```
<bb 2>:  
  a_3 = 1;  
  b_4 = 2;  
  c_5 = 3;  
  n_6 = 6;  
  goto <bb 4>;
```

...

```
<bb 6>:  
  D.1200_8 = a_1 + 2;  
  a_9 = D.1200_8 + 3;
```



Second Level Copy Propagation

Input dump: ccp.c.023t.ccp1

```
<bb 6>:
  D.1200_8 = a_1 + 2;
  a_9 = D.1200_8 + 3;

<bb 7>:
  # a_2 = PHI <a_1(5), a_9(6)>
  a_10 = a_2;
  return a_10;
```

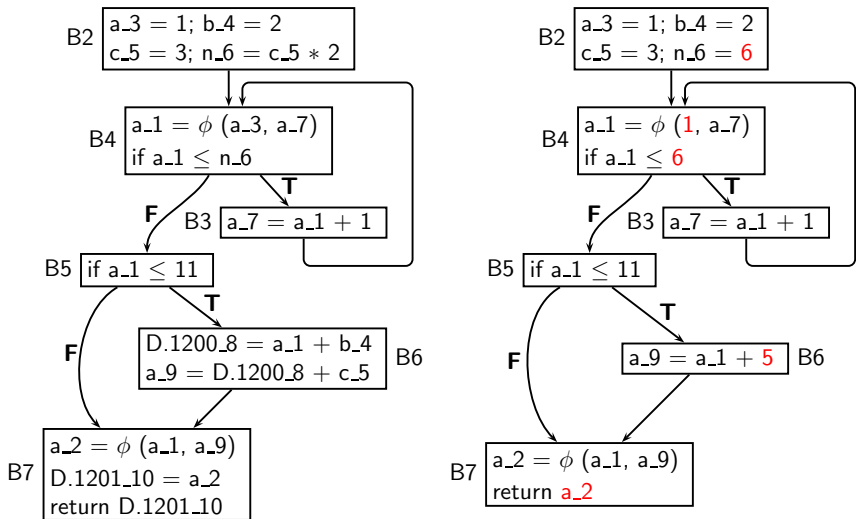
Output dump: ccp.c.027t.copyprop1

```
<bb 6>:
  a_9 = a_1 + 5;

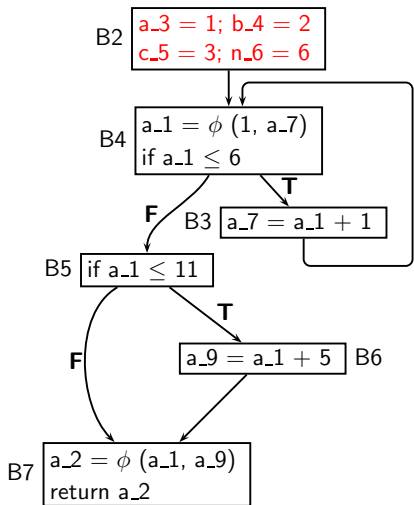
<bb 7>:
  # a_2 = PHI <a_1(5), a_9(6)>
  return a_2;
```



The Result of Copy Propagation and Renaming



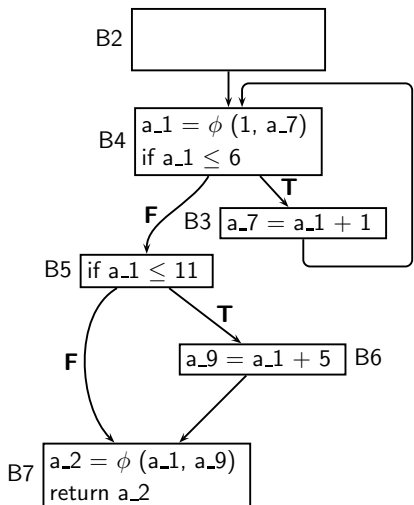
The Result of Copy Propagation and Renaming



- No uses for variables a_3 , b_4 , c_5 , and n_6
- Assignments to these variables can be deleted



Dead Code Elimination Using Control Dependence



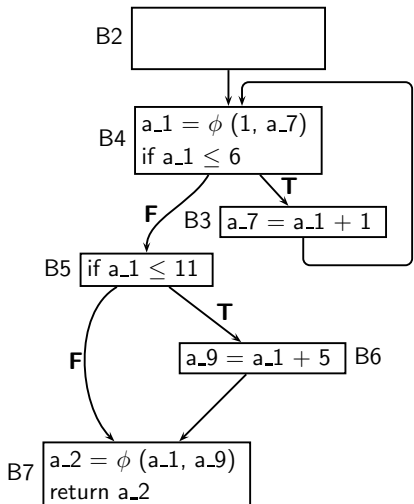
Dump file `ccp.c.029t.cddcel`

```

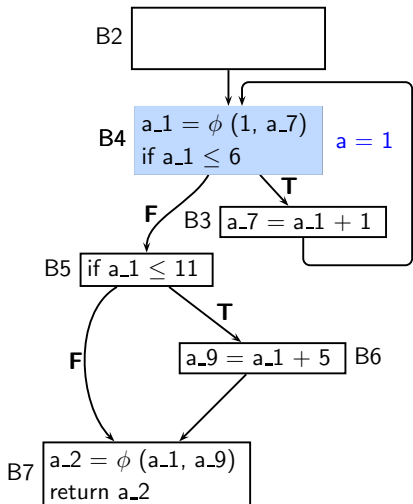
<bb 2>:
    goto <bb 4>;
<bb 3>:
    a_7 = a_1 + 1;
<bb 4>:
    # a_1 = PHI <1(2), a_7(3)>
    if (a_1 <= 6) goto <bb 3>;
    else goto <bb 5>;
<bb 5>:
    if (a_1 <= 11) goto <bb 6>;
    else goto <bb 7>;
<bb 6>:
    a_9 = a_1 + 5;
<bb 7>:
    # a_2 = PHI <a_1(5), a_9(6)>
    return a_2;
  
```



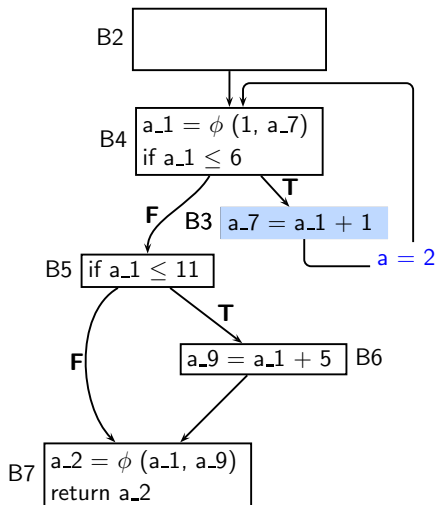
Loop Unrolling



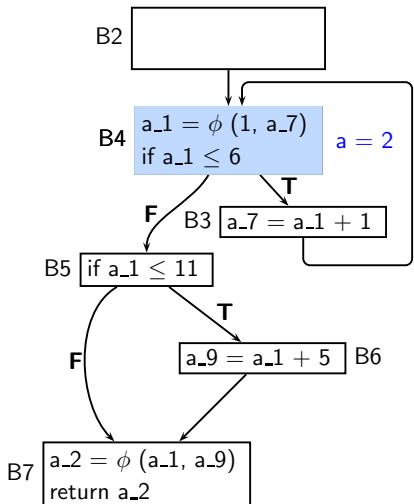
Loop Unrolling



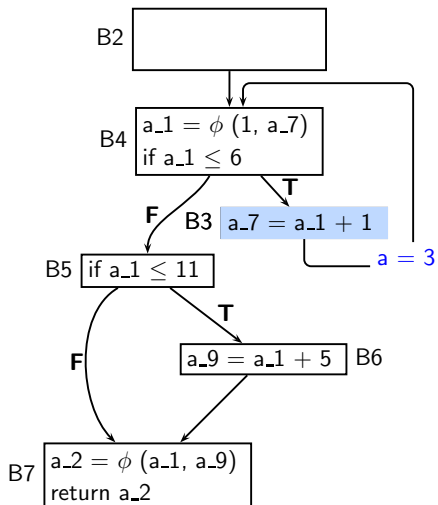
Loop Unrolling



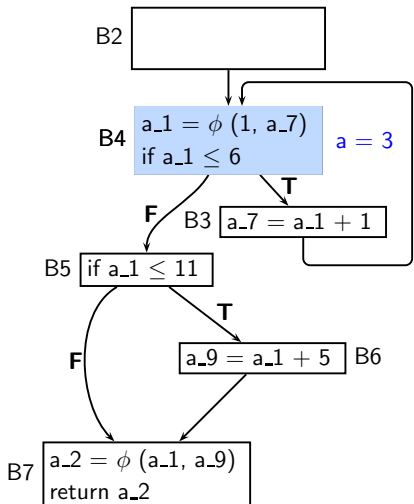
Loop Unrolling



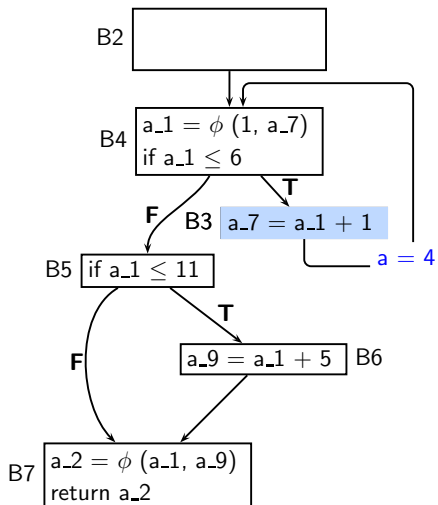
Loop Unrolling



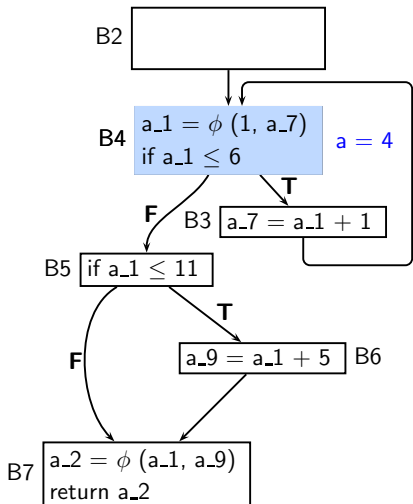
Loop Unrolling



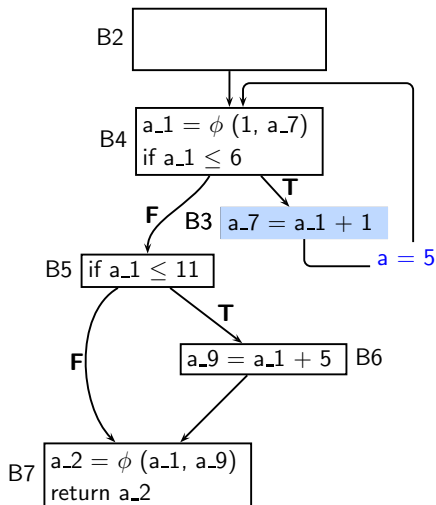
Loop Unrolling



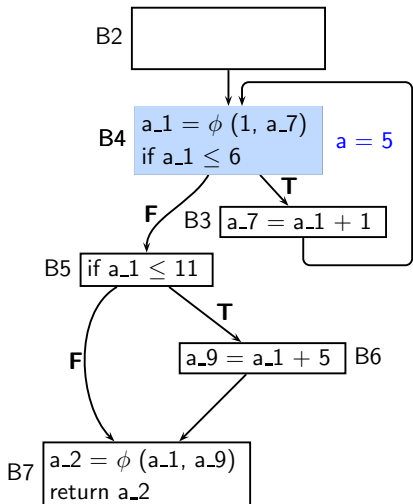
Loop Unrolling



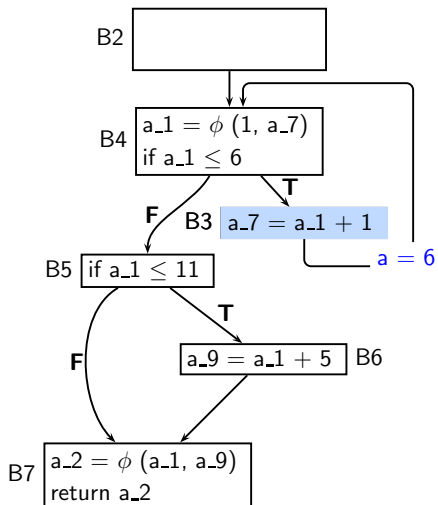
Loop Unrolling



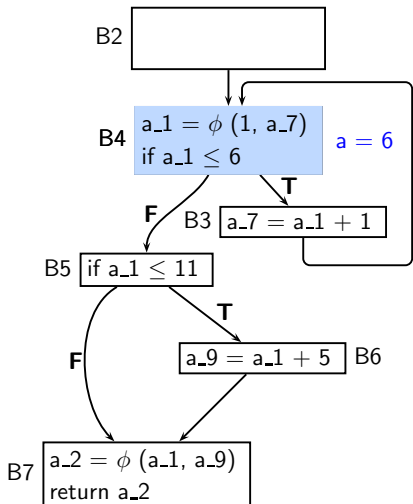
Loop Unrolling



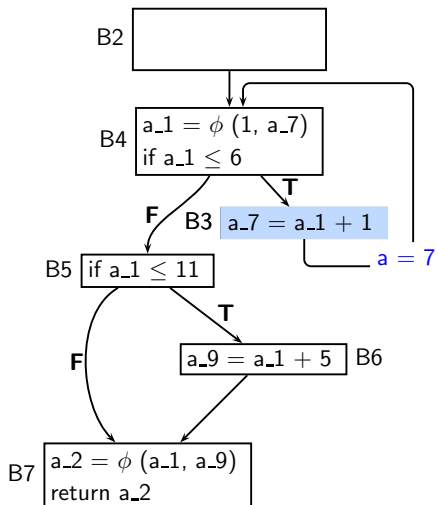
Loop Unrolling



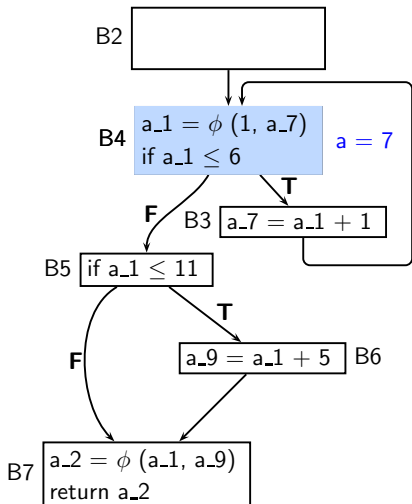
Loop Unrolling



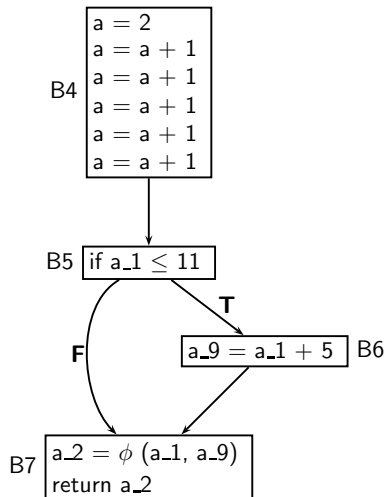
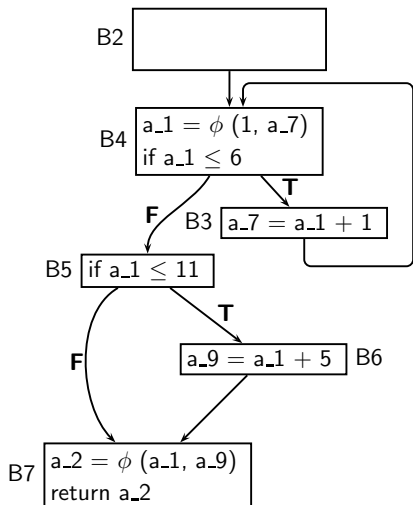
Loop Unrolling



Loop Unrolling



Loop Unrolling



Complete Unrolling of Inner Loops

Dump file: ccp.c.058t.cunrolli

<bb 2>:

```

a_12 = 2;
a_14 = a_12 + 1;
a_16 = a_14 + 1;
a_18 = a_16 + 1;
a_20 = a_18 + 1;
a_22 = a_20 + 1;
if (a_22 <= 11) goto <bb 3>;
else goto <bb 4>;

```

<bb 3>:

```

a_9 = a_22 + 5;

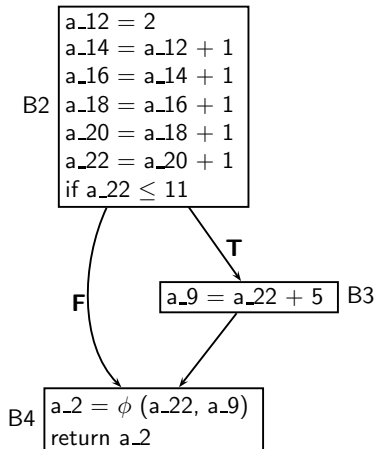
```

<bb 4>:

```

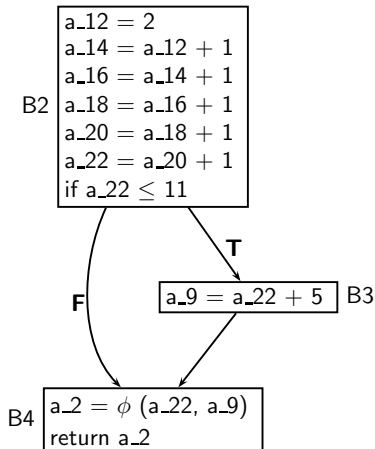
# a_2 = PHI <a_22(2), a_9(3)>
return a_2;

```



Another Round of Constant Propagation

Input



Dump file: ccp.c.059t.ccp2

```

main ()
{
  <bb 2>:
    return 12;
}
      
```



Part 2

Second Example Program

Example Program 2

```
int f(int b, int c, int n)
{ int a;

  do
  {
    a = b+c;
  }
  while (a <= n);

  return a;
}
```

We use this program to illustrate the following optimizations:

Partial Redundancy Elimination,
Copy Propagation, Dead Code
Elimination



Compilation Command

```
$gcc -fdump-tree-all -O2 -S ccp.c
```



Example Program 2

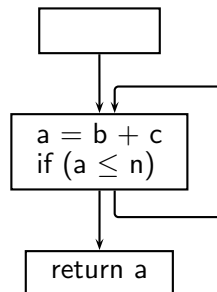
loop.c

```
int f(int b, int c, int n)
{ int a;

  do
  {
    a = b+c;
  }
  while (a <= n);

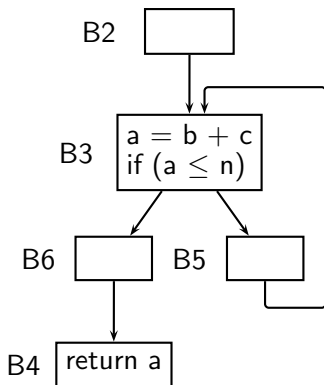
  return a;
}
```

Control Flow Graph



Dump of Input to PRE Pass

Control Flow Graph



loop.c.091t.critcd

<bb 2>:

<bb 3>:

a_3 = b_1(D) + c_2(D);

if (a_3 <= n_4(D)) goto <bb 5>;

else goto <bb 6>;

<bb 5>:

goto <bb 3>;

<bb 6>:

<bb 4>:

a_6 = PHI <a_3(6)>

return a_6;



Input and Output of PRE Pass

loop.c.091t.critcd

<bb 2>:

<bb 3>:

```
a_3 = b_1(D) + c_2(D);  
if (a_3 <= n_4(D))  
    goto <bb 5>;  
else goto <bb 6>;
```

<bb 5>:

```
goto <bb 3>;
```

<bb 6>:

<bb 4>:

```
# a_6 = PHI <a_3(6)>  
return a_6;
```

loop.c.092t.pre

<bb 2>:

```
pretmp.2_7 = b_1(D) + c_2(D);
```

<bb 3>:

```
a_3 = pretmp.2_7;  
if (a_3 <= n_4(D))  
    goto <bb 5>;  
else goto <bb 6>;
```

<bb 5>:

```
goto <bb 3>;
```

<bb 6>:

<bb 4>:

```
# a_6 = PHI <a_3(6)>  
return a_6;
```



Copy Propagation after PRE

loop.c.092t.pre

```
<bb 2>:  
  pretmp.2_7 = b_1(D) + c_2(D);
```

```
<bb 3>:  
  a_3 = pretmp.2_7;  
  if ( a_3 <= n_4(D) )  
    goto <bb 5>;  
  else goto <bb 6>;
```

```
<bb 5>:  
  goto <bb 3>;
```

```
<bb 6>:
```

```
<bb 4>:
```

```
  # a_6 = PHI <a_3(6)>  
  return a_6;
```

loop.c.097t.copyprop4

```
<bb 2>:  
  pretmp.2_7 = b_1(D) + c_2(D);
```

```
<bb 3>:  
  a_3 = pretmp.2_7;  
  if ( n_4(D) >= pretmp.2_7 )  
    goto <bb 4>;  
  else  
    goto <bb 5>;
```

```
<bb 4>:  
  goto <bb 3>;
```

```
<bb 5>:
```

```
  # a_8 = PHI <pretmp.2_7(3)>  
  a_6 = a_8;  
  return a_8;
```



Dead Code Elimination

loop.c.097t.copyprop4

```
<bb 2>:
    pretmp.2_7 = b_1(D) + c_2(D);

<bb 3>:
    a_3 = pretmp.2_7;
    if (n_4(D) >= pretmp.2_7)
        goto <bb 4>;
    else
        goto <bb 5>;

<bb 4>:
    goto <bb 3>;

<bb 5>:
    # a_8 = PHI <pretmp.2_7(3)>
    a_6 = a_8;
    return a_8;
```

loop.c.098t.dceloop1

```
<bb 2>:
    pretmp.2_7 = b_1(D) + c_2(D);

<bb 3>:
    if (n_4(D) >= pretmp.2_7)
        goto <bb 4>;
    else
        goto <bb 5>;

<bb 4>:
    goto <bb 3>;

<bb 5>:
    # a_8 = PHI <pretmp.2_7(3)>
    return a_8;
```



Redundant ϕ Function Elimination and Copy Propagation

loop.c.098t.dceloop1

```
<bb 2>:
  pretmp.2_7 = b_1(D) + c_2(D);

<bb 3>:
  if (n_4(D) >= pretmp.2_7)
    goto <bb 4>;
  else
    goto <bb 5>;

<bb 4>:
  goto <bb 3>;

<bb 5>:
  # a_8 = PHI <pretmp.2_7(3)>
  return a_8;
```

loop.c.125t.phicprop2

```
<bb 2>:
  pretmp.2_7 = c_2(D) + b_1(D);
  if (n_4(D) >= pretmp.2_7)
    goto <bb 4>;
  else
    goto <bb 3>;

<bb 3>:
  return pretmp.2_7;

<bb 4>:
  goto <bb 4>;
```



Final Assembly Program

loop.c.125t.phicprop2

```
<bb 2>:
  pretmp.2_7 = c_2(D) + b_1(D);
  if (n_4(D) >= pretmp.2_7)
    goto <bb 4>;
  else
    goto <bb 3>;

<bb 3>:
  return pretmp.2_7;

<bb 4>:
  goto <bb 4>;
```

loop.s

```
movl    8(%esp), %eax
addl    4(%esp), %eax
cmpl    %eax, 12(%esp)
jge     .L2
rep
ret

.L2:
.L3:
        jmp     .L3
```

Why infinite loop?



Infinite Loop in Example Program 2

```
int f(int b, int c, int n)
{ int a;

  do
  {
    a = b+c;
  }
  while (a <= n);

  return a;
}
```

The program does not terminate
unless $a > n$



Part 3

Conclusions